

Limits on the storage of quantum information in a volume of space

Steven T. Flammia,^{1,2} Jeongwan Haah,^{3,2} Michael J. Kastoryano,⁴ and Isaac H. Kim^{5,6,7}

¹*Centre for Engineered Quantum Systems,*

School of Physics, The University of Sydney, Australia

²*Center for Theoretical Physics, Massachusetts Institute of Technology, Cambridge, USA*

³*Station Q Quantum Architectures and Computation Group,*

*Microsoft Research, Redmond, Washington, USA**

⁴*NBIA, Niels Bohr Institute, University of Copenhagen, Denmark*

⁵*IBM T. J. Watson Research Center, Yorktown Heights, New York, USA*

⁶*Perimeter Institute for Theoretical Physics, Waterloo ON N2L 2Y5, Canada*

⁷*Institute for Quantum Computing, University of Waterloo, Waterloo ON N2L 3G1, Canada*

(Dated: October 20, 2016)

We study the fundamental limits on the reliable storage of quantum information in lattices of qubits by deriving tradeoff bounds for approximate quantum error correcting codes. We introduce a notion of local approximate correctability and code distance, and give a number of equivalent formulations thereof, generalizing various exact error-correction criteria. Our tradeoff bounds relate the number of physical qubits n , the number of encoded qubits k , the code distance d , the accuracy parameter δ that quantifies how well the erasure channel can be reversed, and the locality parameter ℓ that specifies the length scale at which the recovery operation can be done. In a regime where the recovery is successful to accuracy ϵ that is exponentially small in ℓ , which is the case for perturbations of local commuting projector codes, our bound reads $kd^{\frac{2}{D-1}} \leq O(n(\log n)^{\frac{2D}{D-1}})$ for codes on D -dimensional lattices of Euclidean metric. We also find that the code distance of any local approximate code cannot exceed $O(\ell n^{(D-1)/D})$ if $\delta \leq O(\ell n^{-1/D})$. As a corollary of our formulation of correctability in terms of logical operator avoidance, we show that the code distance d and the size \tilde{d} of a minimal region that can support all approximate logical operators satisfies $\tilde{d}d^{\frac{1}{D-1}} \leq O(n\ell^{\frac{D}{D-1}})$, where the logical operators are accurate up to $O((n\delta/d)^{1/2})$ in operator norm. Finally, we prove that for two-dimensional systems if logical operators can be approximated by operators supported on constant-width flexible strings, then the dimension of the code space must be bounded. This supports one of the assumptions of algebraic anyon theories, that there exist only finitely many anyon types.

Quantum information is susceptible to decoherence, but the effect can be mitigated by redundantly encoding the information in a quantum error correcting code. Since a reliable qubit is a scarce resource, it is desirable to achieve maximal protection with a minimum effort, and one of the most promising approaches for achieving this is to incorporate geometric locality into the structure of the code.

Local quantum error correcting codes have been thoroughly studied over the past decade [1], with most of the work focusing on (topological) stabilizer or subsystem codes. A quantum code is said to be *local* if its stabilizer or gauge generators are supported on a geometrically local and bounded region of a lattice embeddable in some metric space. These codes, like most classical and quantum error correcting codes, are often characterized by three numbers $[[n, k, d]]$: n is the number of physical, error-prone qubits comprising the code, k is the maximal number of logical qubits that can be reliably encoded, and d is the distance of the code, i.e. the minimum number of qubits that must be modified to perform a nontrivial logical operation.

Not all values of the triple $[[n, k, d]]$ are achievable, and locality in particular imposes additional constraints. In addition to the constraints imposed by local codes, which we review below, the no-cloning bound implies that exact quantum error-correcting codes cannot correct $n/4$ arbitrary

* jwhaah@microsoft.com

single-qudit errors [2]. In its smallest instance, the bound shows that there is no code on four physical qubits that exactly corrects an arbitrary single-qudit error.

However, the main goal of quantum error correction is to *reduce* effective error rates, and it suffices to only approximately correct errors for this purpose, provided that the approximation error is sufficiently low. Interestingly, the no-cloning bound breaks down when one considers such *approximate* quantum error correcting codes. Leung *et al.* [3] have shown this explicitly by constructing a four-qubit code that approximately corrects an arbitrary single-qubit amplitude damping error. This effect was demonstrated even more dramatically by Crépeau *et al.* [4], who constructed approximate codes that can approximately correct $\lfloor (n-1)/2 \rfloor$ arbitrary single qudit errors, with an approximation error that is exponentially small in n . This demonstrates that approximate codes can radically outperform exact codes by some measures.

Many local quantum codes can be represented as ground spaces of gapped local Hamiltonians. Often these systems have an exact unfrustrated ground space, but there are much more general systems where we expect code properties to hold with some suitable notion of approximation. Known examples include fractional quantum Hall states [5] and Kiteav's honeycomb model in the gapped phase [6]. Might these more general systems allow for unique possibilities for storing quantum information? Many aspects of the unfrustrated Hamiltonian quantum codes remain stable under small perturbations, such as the energy splitting in the ground space and the gap [7]. It is also known that the logical operators, once perturbed, become dressed operators that are quasi-local [8, 9]. However, in light of the examples of Refs. [3, 4], the stability of all code properties for general systems deserves a careful and thorough examination.

Motivated by these observations, we initiate the study of *local approximate quantum error-correcting codes* (local AQEC). It is one of the aims of the present paper to establish a framework allowing the analysis of code properties of subspaces in sufficient generality to encompass many of the interesting gapped many-body models, not necessarily represented by commuting projector Hamiltonians or unfrustrated ground spaces. Our contributions are threefold.

In the first part, we comprehensively analyze the abstract theory of AQEC codes. We introduce a number of different notions of approximate quantum error correction and study the relations between them. In particular, we propose a notion of locally correctable codes for which the class of local commuting projector codes are a subclass. The different notions of correctability that we introduce are all exactly equivalent when the approximation parameter ϵ vanishes, and the relations we establish show that different notions in fact have varying levels of robustness once $\epsilon > 0$.

In the second part, we use these relations to generalize the existing tradeoff bounds for the parameters of local commuting projector codes. These results include the parameter bounds derived by Bravyi, Poulin, and Terhal [10] and the constraints on the structure of the logical operator derived by Haah and Preskill [11] as special cases. Our bounds are only slightly weaker than the prior results for exact codes, implying that local AQEC cannot significantly outperform exact codes by these measures. We furthermore establish a condition under which one of the key axioms of $(2+1)$ -dimensional topological quantum field theory can be derived: If the logical operators can be continuously deformed away from any disk-like region, the ground state degeneracy is bounded by a constant independent of the system size. Roughly, the number of anyon species in 2D is finite.

Our proof techniques make significant departures from much of the previous literature on local codes in that we make heavy use of information-theoretic ideas and methods. We provide information-theoretic generalizations of well known tools from the theory of local codes, including more general forms of the Union Lemma and the Cleaning Lemma. Our proofs and definitions do not require the codes to have local generators; rather all of our analysis is done at the level of subspaces. Locality is only invoked at the level of the recovery operations.

By construction, our results are trivially applicable to local commuting projector codes, but more importantly, every step in our derivations is applicable to the perturbation of such codes.

Thus, we establish that error correction in realistic Hamiltonian systems that only approximate local commuting projector codes is indeed robust to small imperfections.

Prior Work

Previous work on exact quantum codes has placed limits on the achievable code parameters. Beyond the no-cloning bound [2] mentioned above that holds for general exact codes, locality imposes further constraints on the code parameters. Bravyi and Terhal [12] have shown that stabilizer and subsystem codes in D spatial dimensions obey the bound $d \leq O(L^{D-1})$ where $n = L^D$ is the number of physical qubits in the code in a Euclidean lattice. The result on stabilizer codes was extended to the larger class of local commuting projector codes by Bravyi, Poulin, and Terhal [10], who also showed that $kd^{2/(D-1)} \leq O(n)$. Delfosse [13] showed that these arguments can also be adapted to hyperbolic lattices, and showed that surface codes and color codes (with $D = 2$) satisfy $kd^2 \leq O(n(\log n)^2)$, and that this scaling is achievable. The result on subsystem codes was adapted by Bravyi [14] to show that subsystem codes satisfy the bound $kd^{1/(D-1)} \leq O(n)$. Quantum stabilizer codes based on the toric code saturate these bounds for $D = 2$, and subsystem codes that saturate for $D = 2$ [14] or nearly saturate for $D \geq 3$ [15] are also known.

For the class of AQEC, no prior work has shown nontrivial tradeoff bounds like the ones above, and most of the relevant work (such as Ref. [4] above) is confined to codes without any evident locality structure. Known examples of AQEC have generally all been applied to the study of amplitude damping channels (see e.g. [3, 16–18]). Conditions for approximate correction have been derived in terms of the coherent information [19] as well as in terms of error metrics that relate to the Knill-Laflamme conditions for exact error correction [20, 21].

Overview of results

Our results depend on defining appropriate measures of locality and approximation. We will precisely define the relevant measures in Definitions 1 and 2, but, roughly speaking, ℓ will denote a length scale (in lattice spacing units) for the action of a recovery map for erasure errors, and δ will denote the accuracy of the recovery.

Our first contribution is to relate five operationally distinct notions of correctability, decoupling, and cleaning. We will establish similarities between these notions in the approximate recovery setting by deriving inequalities among them without any dimensional factors of the Hilbert spaces involved. When the approximation parameter vanishes, these notions become exactly equivalent.

Decompose the lattice into disjoint regions as $\Lambda = ABC$ and let R denote a purifying system (see e.g. Fig. 1, though there is no restriction on the geometry of the regions). We denote an operator supported on a subsystem using superscripts; for example, ρ^{AB} is a state supported on AB , and ρ^A is its partial trace supported on A . Then we have our first main result.

Result 1 (Imprecise formulation). *The following statements are equivalent.*

- (I) *There exists a local quantum channel \mathcal{R}_B^{AB} with support on AB approximately recovering the erasure of region A : $\mathcal{R}_B^{AB}(\rho^{BCR}) \approx \rho^{ABCR}$.*
- (II) *Tracing out B approximately leaves the system in a product state between A and C : $\rho^{ACR} \approx \omega^A \otimes \rho^{CR}$, for some fixed ω independent of ρ^{ABCR} .*
- (III) *There exists a disentangling unitary U^B on $B = B_1B_2$ that approximately turns any code state into a product state between AB_1 and B_2C : $U^B \rho^{ABCR} U^{B\dagger} \approx \omega^{AB_1} \otimes \rho^{B_2CR}$.*

(IV) *Tracing out B approximately decouples A from C : $\rho^{ACR} \approx \rho^A \otimes \rho^{CR}$.*

(V) *Any logical operator U^{ABC} is associated with another operator V^{BC} supported on the complement of A that act equivalently when restricted to the code space Π : $U^{ABC}|_{\Pi} \approx V^{BC}|_{\Pi}$.*

In particular, we have

$$(I) =_{\ell} (II) =_{\ell} (III) \sim_{\ell} (IV) \sim (V), \quad (1)$$

where ‘ \sim ’ means that the magnitude of the error is not preserved in the implication, and the subscript ‘ ℓ ’ means that the equivalence holds with locality.

Here we have been imprecise about the exact nature of the equivalence, and in particular we have not quantified or defined our notion of approximation or what it means for the equivalence to hold with locality. These statements are given precise formulations and proofs beginning in Section I as follows: (I) Definition 2, (II) Theorem 3, (III) Corollary 4, (IV) Theorem 5, and (V) Theorems 7, 8.

Our next main result shows that for AQECs that have local correctability parametrized by an approximation parameter δ and a length scale ℓ , we have a tradeoff between the capacity and the reliability of encoded information. Consider a Euclidean D -dimensional lattice of linear size L for which we intentionally do not specify any boundary conditions.

Result 2. *For a D -dimensional AQEC with parameters $[[n, k, d, \delta, \ell]]$, it holds that*

$$(1 - c\epsilon \log(1/\epsilon))kd^{\frac{2}{D-1}} \leq c'\ell n^{\frac{2D}{D-1}}, \quad (2)$$

where $\epsilon = n\delta/d$, and $c, c' > 0$ are absolute constants.

In the case that $\delta = 0$ and ℓ is constant, as is the case for local commuting projector codes, we recover the tradeoff bounds of Bravyi, Poulin, and Terhal [10] as a special case. Moreover, we show that for relevant examples (like perturbed versions of local commuting projector codes) the approximation error δ vanishes sufficiently quickly that the inequality remains meaningful in spite of the n dependence inside the parentheses. When the accuracy ϵ is exponentially small in ℓ we obtain the bound quoted in the abstract of $kd^{\frac{2}{D-1}} \leq O(n(\log n)^{\frac{2D}{D-1}})$. This result is restated and proven as Theorem 9 in Section II. Using the approximate cleaning lemma from Section I and similar methods as in the proof of Theorem 9 we prove tradeoff bounds on the support of logical operators of the code (a logical operator is one that preserves the code subspace).

Result 3. *For a D -dimensional AQEC with parameters $[[n, k, d, \delta, \ell]]$ on a lattice of linear size L , if $10L\delta < \ell$, then the code distance is bounded from above by $5\ell L^{D-1}$. Furthermore, there exists a region Y that contains \tilde{d} qubits such that every unitary logical operator U can be approximated by an operator V on Y where*

$$\|(U - V)\Pi\| \leq O(\sqrt{n\delta/d}) \quad \text{and} \quad \tilde{d}d^{\frac{1}{D-1}} \leq O(n\ell^{\frac{D}{D-1}}). \quad (3)$$

Finally, we show how the framework of AQEC can be used to derive one of the the axioms of topological quantum field theory in (2+1) dimensions: the ground space degeneracy is constant independent of system size. We first define a notion of “flexible” logical operator that is made precise in Definition 14, which asserts roughly that logical operators can be deformed across the lattice without changing much how they act on the code space. We then derive the following result.

Result 4. *For any code space Π of a 2-dimensional system admitting flexible logical operators, it holds that*

$$\dim \Pi \leq \exp(c \ell^2), \quad (4)$$

where $c > 0$ is an absolute constant and ℓ is the width of the strip that supports sufficiently faithful logical operators.

The precise formulation and proof of this statement can be found in Section III.

Notation and conventions

As mentioned above, Λ is a regular D dimensional lattice of side length L without any specified boundary conditions. We use letters $A, B, C, D, E, X, Y, Z \in \Lambda$ to denote subsystems. We use Π to denote a subspace of the joint Hilbert space of these subsystems, which we identify as a code space. Unless otherwise specified, different letters stand for disjoint subsystems. We will often say, for example, that a code Π is on ABC when the system into which the code is embedded is divided into three disjoint subsystems. The letter R is reserved for a purifying space of the code space Π . Therefore, as complex vector spaces, Π and the Hilbert space of R are isomorphic. We will use the same symbol Π to denote the projection operator onto the code space. Also, we will write $\rho \in \Pi$ to mean $\rho = \Pi\rho = \rho\Pi$. If a state ρ is pure, we will sometimes use $|\rho\rangle$ to denote a state vector.

When it is necessary to clarify the domain and the codomain of a linear operator, we will use subscripts for domain, and superscripts for codomain. For example, V_B^{AE} denotes a linear operator from B to AE . The subscript will be omitted when the domain is equal to the codomain. For example, V^A is an operator on A . (Under strict practice of our notation, V^A could mean a map from \mathbb{C} to the Hilbert space of A , which would specify a state vector, but we never use such a map.) The same rule applies for quantum channels (i.e. completely positive and trace preserving maps). For example, Tr_A is the erasure channel for system A , mapping from A to scalars. If a channel is tensored with the identity channel, the identity component will be omitted. For instance, we will write $\mathcal{R}_B^{AD}(\sigma^{BC})$, which is a density operator on ACD , in place of $(\mathcal{R}_B^{AD} \otimes \text{id}_C^C)(\sigma^{BC})$. If ρ^{AB} is clear from the context, we will simply write ρ^B in place of $\text{Tr}_A(\rho^{AB})$.

Noise will be modeled by quantum channels \mathcal{N} . It will be convenient to represent the noise in Stinespring form, where $\mathcal{N}(\rho) = \text{Tr}_E(V\rho \otimes \varphi_E V^\dagger)$ with E denoting the purifying system (environment) whose dimension can be taken as the dimension of the channel's input, φ_E is some (pure) state of the purifying system, and V is a unitary operator on the system and purifying space. The complementary channel is obtained by tracing out the system: $\hat{\mathcal{N}}(\rho) = \text{Tr}_S(V(\rho \otimes \varphi_E) V^\dagger)$. We will mainly consider erasure noise. Specifically, \mathcal{N} is chosen to be $\mathcal{N}(\cdot) = \text{Tr}_X(\cdot)$ for some subsystem X .

For any two density operators ρ and σ , the trace distance is defined as

$$\mathfrak{T}(\rho, \sigma) = \frac{1}{2} \|\rho - \sigma\|_1. \quad (5)$$

We will mostly use the one-norm directly instead of the symbol \mathfrak{T} . The fidelity and Bures distance are defined as

$$\mathfrak{F}(\rho, \sigma) = \text{Tr} \sqrt{\sqrt{\sigma} \rho \sqrt{\sigma}} = \max_{|\psi_\rho\rangle, |\psi_\sigma\rangle} |\langle \psi_\rho | \psi_\sigma \rangle|, \quad (6)$$

$$\mathfrak{B}(\rho, \sigma) = \sqrt{1 - \mathfrak{F}(\rho, \sigma)}, \quad (7)$$

where the second equality for the fidelity follows from Uhlmann's theorem [22]. The Bures distance is a metric and satisfies the triangle inequality. These quantities satisfy the Fuchs-van de Graaf relations [23]

$$\mathfrak{F}^2 + \mathfrak{T}^2 \leq 1 \leq \mathfrak{F} + \mathfrak{T}, \quad (8)$$

$$2\mathfrak{B}^2(\rho, \sigma) \leq \|\rho - \sigma\|_1 \leq 2\sqrt{2}\mathfrak{B}(\rho, \sigma). \quad (9)$$

We will often encounter optimizations of the form

$$\max_{\rho^{ABR}} \mathfrak{B}(\rho^{AR}, \rho^A \otimes \rho^R) \quad (10)$$

where a code Π resides on AB . In such an expression, ρ^{ABR} denotes any density matrix of the code vectors possibly entangled with R . Since our discussion will be only for finite dimensional Hilbert spaces, the set of all density matrices is compact under the usual subspace topology inherited from \mathbb{R}^N , which is the metric topology under any mentioned metric. This means that the supremum is always attained at some state.

All log and exp functions have base $e \simeq 2.718$.

I. NOTIONS OF APPROXIMATE QUANTUM ERROR CORRECTION

Most generally, an error correction scheme should be defined in terms of encoding and decoding quantum algorithms [4]. This definition is so general that it includes schemes where a superposition of physical states that are outputs of the encoding algorithm may not correspond to a possible output of the encoding algorithm, e.g. Ref. [4]. In this paper, we restrict ourselves to subspace codes. Namely, we identify a code with a subspace of an n -qubit system, and the encoding map is an isometry from \mathbb{C}^{2^k} to $(\mathbb{C}^2)^{\otimes n}$, where k is the number of encoded qubits.

An exact quantum error correcting code has a distance d if any measurement on fewer than d qubits reveals no information about the encoded state. In other words, a distance d code admits a recovery map that can perfectly reverse the erasure of any $d - 1$ qubits: for any $M \subseteq \Lambda$ such that $|M| < d$, there exists a completely-positive trace-preserving map \mathcal{R} such that

$$\mathcal{R} \circ \text{Tr}_M(\rho) = \rho \quad (11)$$

for all states ρ in the code subspace. An equivalent formulation of the error correction condition, known as the Knill-Laflamme condition [24], expresses Eq. (11) in terms of the Kraus operators of the noise channel.

As remarked earlier, the recovery map will never be perfect in practice and does not have to be perfect; rather, the recovery should be of high enough fidelity to suit the purpose of the code. We use the best recovery map in terms of Bures distance to define approximate quantum error correction.

Definition 1. We say that a region $A \subset \Lambda = AB$ is ϵ -correctable on Π if there exists a recovery operation \mathcal{R}_B^{AB} such that

$$\mathfrak{B}(\mathcal{R}_B^{AB}(\rho^{BR}), \rho^{ABR}) \leq \epsilon, \quad (12)$$

for any code state $\rho^{AB} \in \Pi$ and ρ^{ABR} is any purification. The **code distance** of an approximate error correction code Π is defined as the largest integer d such that any region $A \subseteq \Lambda$ of size $|A| < d$ is ϵ -correctable on Π .

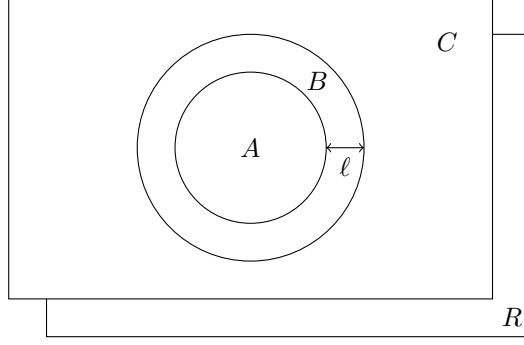


FIG. 1. Decomposition of the lattice in the definitions of local approximate quantum error correction. B shields A from C by a distance at least ℓ , and R represents the purifying space of the code.

The perfect quantum error correction condition corresponds to the $\epsilon = 0$ case.

Our main intuition for local error correcting codes comes from topological systems such as the toric code. There, errors appear as excitation pairs (anyons), that get annihilated when brought together. The recovery of such errors therefore consists of bringing anyon pairs together. In this picture, error recovery is a local operation. This observation motivates the following notion of a locally correctable region for general lattice models.

Definition 2. We say that a region $A \subset \Lambda = ABC$, with $AB = A^{+\ell}$ is the region that includes all qubits within a distance ℓ of A , and C is the complement of AB in Λ , as in Fig. 1, is (δ, ℓ) -correctable on Π if there exists a recovery operation \mathcal{R}_B^{AB} with support on AB such that

$$\mathfrak{B}(\mathcal{R}_B^{AB}(\rho^{BCR}), \rho^{ABCR}) \leq \delta, \quad (13)$$

for any code state $\rho^{ABC} \in \Pi$ and ρ^{ABCR} is any purification.

The parameters of a code can be compactly written as $[[n, k, d, \delta, \ell]]$ for a code on n qubits with the number of encoded qubits k where every region of size less than d is (δ, ℓ) -correctable. By definition, any region that is (δ, ℓ) -correctable is δ -correctable.

One of the central insights of quantum error correction is the information-disturbance tradeoff: the existence of a recovery map implies that the environment knows almost nothing about the state, and vice versa. A very general information disturbance tradeoff was derived in Ref. [20]. See also Refs. [25–27]. Here, we adapt it to our setting by assuming that the noise operation is the partial trace over a given region A . We further sharpen the statement to accommodate for local recovery.

Theorem 3 (Information-Disturbance tradeoff). Let $A \subset \Lambda = ABC$, as in Fig. 1, and define the constant

$$\delta_\ell(A) := \min_{\omega^A} \sup_{\rho^{ABCR}} \mathfrak{B}(\omega^A \otimes \rho^{CR}, \rho^{ACR}), \quad (14)$$

then

$$\inf_{\mathcal{R}_B^{AB}} \sup_{\rho^{ABCR}} \mathfrak{B}(\mathcal{R}_B^{AB}(\rho^{BCR}), \rho^{ABCR}) = \delta_\ell(A), \quad (15)$$

where the inf in Eqn. (15) is over all channels with support on AB .

The proof of Theorem 3 in App. A is quite similar to the ones in Refs. [20, 25].

As a corollary of the information disturbance tradeoff, we show that a region is locally correctable *if and only if* it can be disentangled from its complement by a unitary on the boundary.

Corollary 4. *Let Π be a code space on $ABC = \Lambda$, and R be a purifying auxiliary system. Let $\mathcal{V}_B^{B'B''} = V \cdot V^\dagger$ denote any isometry channel, where B' is some auxiliary system and B'' is a copy of AB . Then,*

$$\delta_\ell(A) = \inf_{\omega^{AB'}, \mathcal{V}_B^{B'B''}} \sup_{\rho^{ABCR}} \mathfrak{B}(\mathcal{V}_B^{B'B''}(\rho^{ABCR}), \omega^{AB'} \otimes \rho^{B''CR}) \quad (16)$$

where $\rho^{B''CR}$ is the same as ρ^{ABCR} but supported on B'' instead of AB .

The proof of the Corollary is in Appendix D.

In Ref. [10], it is shown that for perfect error correcting codes defined by local commuting projectors, there exists a disentangling unitary on the boundary of a correctable region such that it maps any code vector into a product state. Moreover the tensor factor of the product state on the correctable region is the same for any code vector. In view of our expression (16), the result in Ref. [10] corresponds to a situation where B'' and B' are subsystems of B . Although our formulation requires $B'B''$ to have, in general, larger dimension than that of B , our result immediately implies that of Ref. [10] when $\delta_\ell(A) = 0$. This is because for local commuting projector codes, any code state obeys an area law of entanglement Hartley (zeroth Rényi) entropy. This means that there is a unitary transformation that can “compress” the Schmidt components in $B'B''$ of $\omega^{AB'} \otimes \rho^{B''CR}$ into B , yielding a disentangling unitary within B .

The quantity $\delta_\ell(A)$ roughly expresses how weakly a correctable region A is correlated with the far-separated region. Indeed, it is the principle of error correction that the environment, interaction with which would corrupt encoded information, learns nothing about the encoded states. Put differently, the mutual information between a correctable region and the reference system is always small. The quantity $\delta_\ell(A)$ does not directly give the mutual information, but is very similar (proof in Appendix B):

Theorem 5 (Decoupling-Correctibility). *Let $A \subset \Lambda = ABC$ as in Fig. 1, then*

$$\frac{1}{9}\delta_\ell(A)^2 \leq \sup_{\rho^{ABCR}} \mathfrak{B}(\rho^{ACR}, \rho^A \otimes \rho^{CR}) \leq 2\delta_\ell(A) \quad (17)$$

Schumacher and Nielsen [28] defined the coherent information, which is equivalent to the mutual information $I(A : R)$ between the purifying system R and the correctable region A . They argued that if the coherent information does not degrade, then perfect error correction is possible. Schumacher and Westmoreland [19] generalized this to an approximate setting, and argued that the change in the coherent information bounds the recovery fidelity. However, in their argument the recovery map was constructed from a code state, and was not shown to be applicable for all code states. This is not satisfactory for error correction since any correcting map must not know about the code state. In the perfect error correction setting, this is no longer an issue due to the Knill-Laflamme condition [24]. (See also [29].) We were unable to find a reference that proves the existence of a recovery map that would work for all code vectors in the approximate error correction setting, so we proved our own in Theorem 5. Schumacher and Westmoreland’s claim follows qualitatively with a recovery map that depends only on Π because [30, 31]

$$\begin{aligned} I_\rho(A : R) &\geq -2 \log \mathfrak{F}(\rho^{AR}, \rho^A \otimes \rho^R) \\ &\geq 2\mathfrak{B}^2(\rho^{AR}, \rho^A \otimes \rho^R). \end{aligned} \quad (18)$$

Note that the continuity of the mutual information (see Appendix F) implies

$$I_\rho(A : R) \leq 18\sqrt{2} \delta_\ell(A) \log \left(\frac{\dim \Pi}{2\sqrt{2} \delta_\ell(A)} \right). \quad (19)$$

No attempt has been made to optimize this inequality, and in fact we will only use a weaker statement with $\log \dim \Pi = k \log 2 \geq 1$ and $\delta_\ell(A) \leq 1/e$,

$$I_\rho(A : R) \leq O(k \delta_\ell(A) \log(1/\delta_\ell(A))) \quad (20)$$

since some of the formulas become simpler.

A. Cleanability as an alternative notion of correctability

Here we characterize the correctability of a region by logical operator avoidance. If all logical operations can be done outside a region A , then, tautologically, for any logical operator U one can achieve the equivalent action on the code space without touching A . Thus, the region A is *cleaned* of U . In a perfect error correction setting, it is known that a correctable region can be cleaned of any logical operator, and hence the name “cleaning lemma” [11, 12, 32]. We prove a generalization of the cleaning lemma in the approximate setting, and complement it with the converse statement. This establishes another error correction criterion based on the support of logical operators.

Definition 6. An operator U is **logical** if it commutes with the code space projector Π .

Theorem 7 (A correctable region avoids logical operators.). *Suppose A is a (δ, ℓ) -correctable region of the lattice $ABC = \Lambda$; i.e., there exists \mathcal{R}_B^{AB} such that*

$$\sup_{\rho^{ABC}} \mathfrak{B}(\mathcal{R}_B^{AB}(\rho^{BCR}), \rho^{ABCR}) \leq \delta \quad (21)$$

where AB is the ℓ -neighborhood of A . Then, for any logical unitary U^{ABC} , the pull-back $V^{BC} = (\mathcal{R}_B^{AB})^*(U^{ABC})$ satisfies

$$\begin{aligned} \|(U^{ABC} - V^{BC})\Pi\| &\leq 4\sqrt{\delta}, \\ \|\Pi(U^{ABC} - V^{BC})\| &\leq 4\sqrt{\delta}. \end{aligned} \quad (22)$$

The converse without the locality of the recovery map is also true:

Theorem 8 (A region avoiding logical operators is correctable.). *Given the lattice $AB = \Lambda$, suppose for any logical unitary U^{AB} there exists an operator $\|V^B\| \leq 1$ supported on B such that*

$$\begin{aligned} \|(U^{AB} - V^B)\Pi\| &\leq \delta, \\ \|\Pi(U^{AB} - V^B)\| &\leq \delta. \end{aligned} \quad (23)$$

Then there exists ω^A such that

$$\sup_{\rho^{ABR}} \|\rho^{AR} - \omega^A \otimes \rho^R\|_1 \leq 5\delta, \quad (24)$$

and A is $\sqrt{5\delta/2}$ -correctable.

Theorems 7 and 8 are proved in Appendix C.

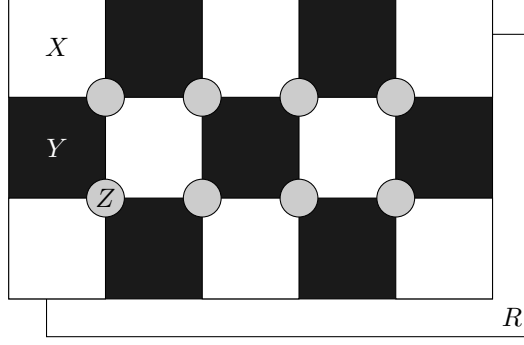


FIG. 2. Decomposition of the lattice in 2D leading to the tradeoff bound of Theorem 9.

II. TRADEOFF BOUNDS FOR LOCALLY CORRECTABLE CODES

Our first main result is a tradeoff bound between the number of encoded qubits and the distance of a code on the lattice. The proof is similar in essence to the tradeoff bound proved for local commuting projector codes by Bravyi, Poulin, and Terhal [10]. However, the details vary in several respects because we are dealing with approximate error correcting codes rather than exact ones. Especially, Ref. [10] uses an algebraic decomposition of the Hilbert space resulting from the representation theory of commuting operators. In the approximate setting we cannot use such a decomposition, but instead we use a decoupling characterization of locally correctable codes. We state and prove the theorem for a two-dimensional Euclidean lattice of qubits for clarity of presentation, but the proof generalizes easily to D -dimensional lattices with $D \geq 2$. Having a Euclidean geometry will be important in the final statement because we are going to use the fact that the surface area of a ball of radius r grows like r^{D-1} . We note that the technique below, which is borrowed from Ref. [10], will be applicable for other geometries.

Theorem 9. *Let Π be a local error correcting code with parameters $[[n, k, d, \delta, \ell]]$ defined on a D -dimensional Euclidean lattice of qubits. It holds that*

$$\left(1 - c \frac{n\delta}{d} \log \frac{d}{n\delta}\right) kd^{\frac{2}{D-1}} \leq c'n\ell^{\frac{2D}{D-1}}. \quad (25)$$

where $c, c' > 0$ are constants independent of any parameters n, k, d, δ, ℓ .

Obviously, this is meaningful only if the factor in the parenthesis is positive; we must have sufficiently small δ . For local commuting projector codes, $\delta = 0$ for ℓ such that every local projector is contained in $\ell \times \ell$ square. We then recover the bound $kd^{2/(D-1)} \leq O(n)$ of Ref. [10]. Note that we did not need to invoke any form of local generators of the code, rather we based our proof exclusively on the natural locality structure of the subspace. If $\delta = \delta(\ell) \sim \exp(-\ell/\xi)$ (as we later show is the case for perturbations of commuting projector codes), then we obtain

$$kd^{\frac{2}{D-1}} \leq O\left(n(\log n)^{\frac{2D}{D-1}}\right) \quad (26)$$

by choosing $\ell \sim \xi \log n \sim \xi \log(L)$, where $L = n^{1/D}$ the linear system size.

The proof largely consists of three steps. We divide the whole lattice into three subsystems X, Y, Z , each of which is a collection of disjoint correctable simply connected regions. They are depicted in Fig. 2 for $D = 2$. The first step is to choose the disconnected components of X and Y as large as possible. The second step is to show that the union X or Y of those components is still

correctable. For this, each component should be sufficiently separated. In both steps, the locality of the error correcting map will be important. The third step is to employ a technique in a proof of the quantum Singleton bound [33] to bound the code space dimension by the volume of Z .

We will need two lemmas below, which allow us to construct regular correctable regions with weight larger than the distance. The *expansion lemma* gives a prescription on how to grow a correctable region, and the *union lemma* shows that the union of two distant correctable regions is again correctable.

Lemma 10 (Expansion Lemma). *Let $AB = A^{+\ell}$ be the ℓ -neighborhood of a region A . If A is (ϵ_A, ℓ) -correctable and B is (ϵ_B, ℓ) -correctable, then $A \cup B$ is $(\epsilon_A + \epsilon_B, \ell)$ -correctable.*

Proof. Let $ABC = A^{+2\ell}$, and let D be the rest of the lattice. Fix ω^A that saturates the optimum for $\delta_\ell(A)$. From the correctability of region B , there exists a local recovery map \mathcal{R}_{AC}^{ABC} such that

$$\mathfrak{B}(\rho^{ABCD R}, \mathcal{R}_{AC}^{ABC}(\rho^{ACDR})) \leq \epsilon_B \quad (27)$$

for any code state $\rho^{ABCD R}$. Define a channel \mathcal{R}_C^{ABC} by

$$\mathcal{R}_C^{ABC}(\sigma^C) := \mathcal{R}_{AC}^{ABC}(\omega^A \otimes \sigma^C) \quad (28)$$

for any state σ^C . Now, for an arbitrary code state $\rho^{ABCD R}$

$$\begin{aligned} & \mathfrak{B}(\mathcal{R}_C^{ABC}(\rho^{CDR}), \rho^{ABCD R}) \\ &= \mathfrak{B}(\mathcal{R}_{AC}^{ABC}(\omega^A \otimes \rho^{CDR}), \rho^{ABCD R}) \\ &\leq \mathfrak{B}(\mathcal{R}_{AC}^{ABC}(\omega^A \otimes \rho^{CDR}), \mathcal{R}_{AC}^{ABC}(\rho^{ACDR})) \\ &\quad + \mathfrak{B}(\mathcal{R}_{AC}^{ABC}(\rho^{ACDR}), \rho^{ABCD R}) \\ &\leq \mathfrak{B}(\omega^A \otimes \rho^{CDR}, \rho^{ACDR}) + \epsilon_B \\ &\leq \epsilon_A + \epsilon_B, \end{aligned} \quad (29)$$

where we used monotonicity of the Bures distance. Therefore, the map \mathcal{R}_C^{ABC} recovers AB from C up to error $\epsilon_A + \epsilon_B$. \square

Lemma 11 (Union Lemma). *Suppose two regions A and B are separated by distance at least ℓ , where A is (ϵ_A, ℓ) -correctable and B is ϵ_B -correctable. Then the union AB is $(\epsilon_A + \epsilon_B)$ -correctable.*

Proof. Fix ω^A and ω^B that saturate the optima for $\delta_\ell(A)$ and $\delta_\ell(B)$. The claim is proved as

$$\begin{aligned} \mathfrak{B}(\rho^{ABR}, \omega^A \omega^B \rho^R) &\leq \mathfrak{B}(\rho^{ABR}, \omega^A \rho^{BR}) \\ &\quad + \mathfrak{B}(\omega^A \rho^{BR}, \omega^A \omega^B \rho^R) \\ &\leq \epsilon_A + \mathfrak{B}(\rho^{BR}, \omega^B \rho^R) \end{aligned} \quad (30)$$

$$\leq \epsilon_A + \epsilon_B \quad (31)$$

for any code state ρ^{ABCR} . \square

Proof of Theorem 9. The statement becomes vacuous if $d \leq \ell$. So, assume $d > \ell$.

(Step 1) Since $d > \ell$, a single site is (δ, ℓ) -correctable. Its boundary contains $O(\ell^D)$ qubits. If this is less than d , then it is also (δ, ℓ) -correctable, and by Lemma 10 an enlarged hypercube of linear size of order ℓ is $(2\delta, \ell)$ -correctable. After m steps of induction, we obtain a hypercube of linear size of order $m\ell$ that is $(O(m\delta), \ell)$ -correctable. The induction must stop if the hypercube is so large that the boundary “area” $O(m^{D-1}\ell^D)$ becomes larger than d . Thus, m can be as large

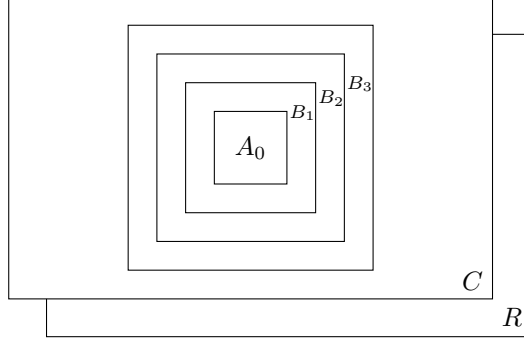


FIG. 3. Construction of the largest correctable square by successively adding rings of size at most d .

as $O\left((d/\ell^D)^{\frac{1}{D-1}}\right)$. Therefore, we conclude that any hypercube of linear size $O\left((d/\ell)^{\frac{1}{D-1}}\right)$ is $\left(O\left(\delta(d/\ell^D)^{\frac{1}{D-1}}\right), \ell\right)$ -correctable.

(Step 2) Consider a decomposition of the lattice into three regions $XYZ = \Lambda$, where X and Y are, respectively, the unions of disconnected hypercubes of linear size $O\left((d/\ell)^{\frac{1}{D-1}}\right)$ constructed in Step 1 with their $(D-2)$ -dimensional “corners” of width ℓ removed. The hypercubes of X and Y are arranged in a high-dimensional checkerboard. Z is the rest of the lattice. See Fig. 2 for an illustration of the decomposition for $D = 2$. Let L be the linear system size. The collections X and Y consist of $O\left(L(\ell/d)^{\frac{1}{D-1}}\right)^D$ hypercubes, respectively. Applying Lemma 11 inductively over all squares of X , we see that X is ϵ -correctable where

$$\epsilon = O(1) \left(L(\ell/d)^{\frac{1}{D-1}}\right)^D \cdot \left(\delta(d/\ell^D)^{\frac{1}{D-1}}\right) = O(1) \frac{n\delta}{d}. \quad (32)$$

Similarly, Y is also ϵ -correctable.

(Step 3) Let $\rho^{XYZ} = \Pi / \dim \Pi$ be the maximally mixed code state, and ρ^{XYZR} be a purification. By definition, $S(\rho^R) = \log \dim \Pi = k \log 2$. Since X and Y are ϵ -correctable, respectively, by Eq. (20), the mutual information to R must be small:

$$S(\rho^X) + S(\rho^R) - S(\rho^{XR}) \leq O(1)k\epsilon \log(1/\epsilon), \quad (33)$$

$$S(\rho^Y) + S(\rho^R) - S(\rho^{YR}) \leq O(1)k\epsilon \log(1/\epsilon). \quad (34)$$

Adding the two inequalities and using the fact that ρ^{XYZR} is pure, we obtain

$$S(\rho^X) + S(\rho^Y) + 2k \log 2 - S(\rho^{YZ}) - S(\rho^{XZ}) \leq O(1)k\epsilon \log(1/\epsilon).$$

By the subadditivity of entropy (where \tilde{O} hides log factors), it follows that

$$k(1 - \tilde{O}(\epsilon)) \leq O(1)S(\rho^Z) \leq O(|Z|). \quad (35)$$

The region Z consists of a grid of “bar segments”, separating the individual hypercubes in X and Y . A bar segment has two sides of length ℓ and $D-2$ sides of length $O(d/\ell)^{1/(D-1)}$. There are as many bar segments as there are the hypercubes of X and Y , the number of which is $O(n)(\ell/d)^{D/(D-1)}$. Therefore,

$$k(1 - \tilde{O}(\epsilon)) \leq O(n)(\ell/d)^{\frac{D}{D-1}} \cdot \ell^2 \cdot O(d/\ell)^{\frac{D-2}{D-1}} = O(n)\ell^{\frac{2D}{D-1}}d^{\frac{-2}{D-1}}. \quad (36)$$

This complete the proof of Theorem 9. \square

A. Support of Logical Operators

As corollaries of the preceding results, we can also derive constraints on the support of logical operators.

Bravyi and Terhal [12] have shown that for stabilizer codes defined by local stabilizer generators on a Euclidean lattice there always exists a nontrivial logical operator that is supported on a thin slab $((D - 1)$ -dimensional). This result was generalized to local commuting projector codes [11]. In two-dimensional lattices, this suggests that under generic interaction of the system with an environment there would be a process subject to a constant energy penalty that implements a nontrivial logical operation, i.e., an error, on the encoded state. This is a strong argument against *self-correction* in two-dimensional systems under a thermalizing interaction where the code space is the ground space. In other words, either bit-flip errors or phase errors would occur by thermalization on an encoded qubit at a rate independent of code distance.

However, this does not immediately rule out the possibility that only one of bit or phase information is corrupted by thermalization, but the other information is protected. Indeed, any ferromagnetic (symmetry broken) system is protected from bit-flip errors. For this observation, Haah and Preskill [11] have asked whether it is possible to have a *partially self-correcting* quantum memory in two-dimensions where the code distance is high, but, e.g., bit-flip errors are suppressed under thermalization. It was found that in any two-dimensional local commuting projectors code if the code distance d , then *all* logical operators can be supported on a region that contains $\tilde{d} = O(L^2/d)$ physical qubits. This is a negative result towards partial self-correction because having a large distance $d \sim L$ implies that any other logical operator lives on a network of finitely many string-like regions.

We generalize these results using the machinery we have developed.

Theorem 12. *For any (δ, ℓ) -correctable code Π with $\dim \Pi > 1$ on a D -dimensional lattice of linear size L , if $10L\delta < \ell$, then the code distance is bounded from above by $5\ell L^{D-1}$.*

Proof. The argument is essentially one-dimensional, and it suffices to prove the theorem for $D = 1$. Suppose on the contrary to the claim that the any segment of length 5ℓ is (δ, ℓ) -correctable. In particular, a segment of length 5ℓ is (δ, ℓ) -correctable, and its two-component boundary B of size 2ℓ is also (δ, ℓ) -correctable. The expansion lemma 10 implies that the union AB of length 7ℓ is $(2\delta, \ell)$ -correctable. The two-component boundary of AB is also (δ, ℓ) -correctable, and we can again apply the expansion lemma 10 to have $(3\delta, \ell)$ -correctable region of length 9ℓ . After $O(L/\ell)$ times of iteration, we see that entire system is $(\delta L/\ell)$ -correctable. This is a contradiction since the entire system is certainly not ϵ -correctable with $\epsilon < 1$. \square

This code distance bound is intimately related to the absence of topological order (without any symmetry) in one dimension $D = 1$. If we regard the ground space as a code space, then having a small code distance means that there is an operator of small support that takes different expectation values for distinct ground states, indicating that the degeneracy will be lifted upon perturbation by that operator. In particular, if the ground space is locally correctable, which is the case if the quantum phase of the one-dimensional system is represented by a commuting Hamiltonian by the results in Section IV below, then the ground space degeneracy, if any, is lifted by perturbations. For fermionic systems, this only says the ground space is not locally correctable, or a degeneracy-lifting local operator of fermion parity *even or odd* should exist.

Theorem 13. *For any (δ, ℓ) -correctable code of code distance d on a D -dimensional lattice with Euclidean geometry of linear size L , there exists a region Y that contains \tilde{d} qubits such that every*

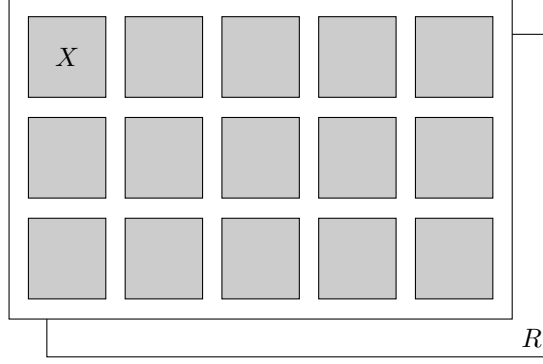


FIG. 4. Decomposition of the lattice used for Theorem 13.

logical operator U can be approximated by an operator V on Y where

$$\begin{aligned} \|(U - V)\Pi\| &\leq O\left(\sqrt{n\delta/d}\right) \\ \tilde{d}d^{\frac{1}{D-1}} &\leq O(n\ell^{\frac{D}{D-1}}) \end{aligned} \quad (37)$$

If δ becomes exactly zero at some ℓ , independent of L , then the conclusion becomes that of Ref. [11].

Proof. We divide the system as in Fig. 4, where each hypercube is separated by a distance at least ℓ . Each hypercube in X has linear size at most $O(d/\ell)^{\frac{1}{D-1}}$, and is at least $\left(O(\delta(d/\ell^D)^{\frac{1}{D-1}}), \ell\right)$ -correctable by Step 1 in the proof of Theorem 9. The union X of all such squares is $O(n\delta/d)$ -correctable by the union lemma 11. (See Eq. (32)). Then, Theorem 7 (cleaning lemma) implies that the complement, on which there are $O(n)(\ell/d)^{\frac{D}{D-1}} \cdot \ell \cdot O(d/\ell)$ qubits, supports all logical operators to accuracy $O(\sqrt{n\delta/d})$ in operator norm. \square

III. FLEXIBLE STRINGS IMPLY FINITE DEGENERACY

Here we apply Theorem 8 to topologically ordered two-dimensional systems, to show that under a mild condition the degeneracy can be at most a constant independent of system size. This constant degeneracy holds for all examples we know of with a stable ground state subspace subject to arbitrary perturbations, and is intimately related to one of the core assumptions of algebraic anyon theories (modular tensor categories) that there are only finitely many superselection sectors (simple objects). Without the finiteness, any interesting computation in the algebraic anyon theory would contain an infinite sum, rendering the theory substantially different from what we understand with the finiteness. For instance Vafa's theorem [34] stating that the topological spin is rational would not hold without the finiteness. Also, this is a unique property of two dimensions, since in higher dimensions it is trivially possible to have very large degeneracy that is stable under perturbations.

We consider periodic boundary conditions, though it will not be too important whether the boundary conditions are open or periodic. The overall topology of the system is a 2-torus.

In a topologically ordered two-dimensional system, it is expected that

- (o) there exists a complete set of operators acting within the ground space Π (logical operators) such that they are supported along a narrow strip (string) wrapping around non-contractible loops of the system.

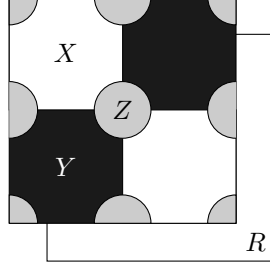


FIG. 5. Decomposition of the lattice in 2D used to define flexible logical operators. The diameter of the disk Z is ℓ .

It is also expected that the string can be bent while implementing the same action on the ground space as long as the support remains isotopic to the initial one. These properties can be regarded as the mathematical defining properties of topological order.

Assuming the existence of a complete set of such operators acting within Π , we now bound the degeneracy. Consider four squares with corners removed that fills the system as in Fig. 5. The corners are occupied by four regions comprising Z . The two squares on one diagonal comprise X , and the other two squares comprise Y . The size of a disk in Z has diameter ℓ so that the two squares in either X or Y are separated by distance ℓ . According to the property (o), we can find a complete set of logical operators supported on XZ since XZ contains two essential non-contractible loops of the torus. It may be unreasonable to assume that all operators that strictly preserve the ground space are supported entirely on XZ . Instead, we may reasonably assume that such an operator can be approximated by one that is supported entirely on XZ , and the approximation becomes better as we increase ℓ .

Hence, we are well motivated to define the notion of flexible logical operators as follows:

Definition 14. *A subspace Π on a two-dimensional system admits flexible (logical) operators if for any logical unitary operator U^{XYZ} there exist operators V_1^{YZ} supported on YZ and V_2^{XZ} on XZ such that $\|V_i\| \leq 1$, $\|\Pi(U - V_i)\| \leq \epsilon_\ell$, and $\|(U - V_i)\Pi\| \leq \epsilon_\ell$, where $i = 1, 2$ and ϵ_ℓ is independent of system size and vanishes as $\ell \rightarrow \infty$.*

From this definition Theorem 8 asserts that the complement $Y = (XZ)^c$ is $\sqrt{5\epsilon_\ell/2}$ -correctable. Interchanging X and Y , we conclude that Y is also $\sqrt{5\epsilon_\ell/2}$ -correctable. Next, we employ the technique of Step 3 in the proof of Theorem 9. Recalling the continuity of mutual information Eq. (20), we deduce that

$$\left(1 - \tilde{O}(\epsilon_\ell^{1/2})\right) \log \dim \Pi \leq S(\rho^Z) \leq \ell^2 \log 2. \quad (38)$$

Choosing ℓ sufficiently large (and hence ϵ_ℓ sufficiently small), we have

$$\dim \Pi \leq \exp(O(\ell^2)). \quad (39)$$

In summary, the bound on the degeneracy $\dim \Pi$ is determined by the width of the strip that can support sufficiently faithful logical operators. It is noteworthy that the error from restricting the logical operator onto the strip only has to be suppressed to an absolute constant. Also, the exponent ℓ^2 being quadratic in ℓ is optimal; ℓ^2 copies of the toric code can be laid on a lattice to achieve this quadratic scaling.

IV. ROBUSTNESS OF LOCAL COMMUTING PROJECTOR CODES

The results so far use properties of subspaces, and do not hinge on particular code constructions. Here we will explicitly construct a class of local approximate error correction codes, as the perturbation of exact codes. Let $H_0 = -\sum_{k \in \Lambda} h_k$ be a local unfrustrated commuting projector Hamiltonian. That is, each term is a projector $h_k = h_k^2$, commutes with any other term $[h_k, h_j] = 0$, and is supported on a region around site k of diameter less than w (the interaction range), and any ground state $|\psi\rangle$ satisfies $h_k |\psi\rangle = |\psi\rangle$. The code subspace is identified with the ground state subspace of H_0 : $\Pi = \prod_{k \in \Lambda} h_k$. We summarize the error correction properties of these *local commuting projector* codes in the following lemma:

Lemma 15. *Let Π be a commuting projector code, and $ABC = \Lambda$ be decomposition of the lattice such that the distance between A and C is at least $\ell \geq w$, the interaction range (e.g. as in Fig. 3.) Then the following are equivalent:*

- (i) *Topological Quantum Order (TQO): for any observable O_A with support on A , any two ground states $|\phi\rangle$ and $|\psi\rangle$ give the same expectation value, $\langle\phi| O^A |\phi\rangle = \langle\psi| O^A |\psi\rangle$.*
- (ii) *Decoupling: For any $\rho \in \Pi$, we have $I_\rho(A : CR) = 0$.*
- (iii) *Error correction: There exists a recovery map acting on AB such that $\mathcal{R}_B^{AB}(\rho^{BC}) = \rho^{ABC}$ for any $\rho \in \Pi$.*
- (iv) *Disentangling unitary: For any $\rho \in \Pi$, there exists a unitary U^B , such that $U^B \rho U^{B\dagger} = \omega^{AB_1} \otimes \rho^{B_2C}$, for some state ω^{AB_1} .*
- (v) *Cleaning: For any unitary U preserving the code space, there exists a unitary V^{BC} such that $U|_\Pi = V^{BC}|_\Pi$.*

The condition (i) is Knill-Laflamme [24] condition for correctability. The implication from (iii) to (iv) was derived in Ref. [10] with slightly larger B . The implication from (iv) to (v) was derived in Ref. [11] with a slightly larger B .

Proof. (i) \Rightarrow (ii): (i) is equivalent to $\Pi O^A \Pi = c(O^A) \Pi$ for some scalar $c(O^A)$. Then, it suffices to look at the correlation function between A and C ;

$$\langle\psi| O^A O^C |\psi\rangle \tag{40}$$

$$= \langle\psi| \Pi O^A \left(\prod_{k: \ell\text{-away from } A} h_k \right) \left(\prod_{k: \ell\text{-away from } C} h_k \right) O^C \Pi |\psi\rangle \tag{41}$$

$$= \langle\psi| \Pi O^A \Pi O^C \Pi |\psi\rangle \tag{42}$$

$$= c(O^A) c(O^C) = \langle\psi| O^A |\psi\rangle \langle\psi| O^C |\psi\rangle. \tag{43}$$

(ii) \Leftrightarrow (iii) \Leftrightarrow (iv) by Theorem 3, Corollary 4, and Theorem 5 with $\delta_\ell(A) = 0$.

(iii) \Rightarrow (v) by Theorem 7 with $\delta = 0$.

(v) \Rightarrow (i): Theorem 8 with $\delta = 0$ implies that the reduced density matrix for A is the same for all code states. \square

We now consider weak perturbations of commuting projector codes. We assume that the code distance (Definition 1) grows with the system size by a power law ($d = \Omega(n^\gamma)$ for some $\gamma > 0$). In addition, we assume that the Hamiltonian H_0 obeys the “local TQO” condition of Refs. [7, 9].

These conditions ensure that a perturbed Hamiltonian $H_1 = H_0 + \sum_k V_k$ has a *gapped energy spectrum* above the ground state subspace, where the V_k are local and $\|V_k\| \leq \epsilon$ for a constant ϵ that only depends on the spatial dimension and the interaction range w , but the V_k are otherwise arbitrary.

The local TQO condition is similar to the TQO condition, but eliminates effects from the correctable region's boundary. H_0 is said to obey local TQO if

$$\Pi^{A^w} O^A \Pi^{A^w} = c(O^A) \Pi^{A^w} \quad (44)$$

$$\text{where } \Pi^{A^w} = \prod_{k: \text{dist}(k, A) \leq w} h_k \quad (45)$$

for any region A of size less than d . In other words, if any state $|\psi\rangle$ satisfies $h_k |\psi\rangle = |\psi\rangle$ for h_k around a correctable region A , then the reduced density matrix for A of $|\psi\rangle$ is determined and unique.

The gap stability result implies that the ground state subspaces Π_0 and Π_1 of the unperturbed Hamiltonian H_0 and the perturbed Hamiltonian H_1 , respectively, are connected by a locality-preserving unitary U [8, 35, 36].

$$\Pi_1 = U \Pi_0 U^\dagger \quad (46)$$

$$\left\| U O^X U^\dagger - V^{X+r} O^X (V^{X+r})^\dagger \right\| \leq c_1 \|O^X\| \exp(-c_2 r / \log^2 r). \quad (47)$$

Here, X is an arbitrary region, and V^{X+r} is strictly supported on the r -neighborhood of X , and $c_1, c_2 > 0$ are constants. The particular form of the function on the right-hand side of (47) is from Ref. [36, Theorem 3.4].

We are now in position to state the main theorem of this section:

Theorem 16. *Let H be a local commuting frustration-free Hamiltonian, whose ground state defines an $[[n, k, d]]$ quantum error correcting code with $d = \Omega(n^\gamma)$ for some $\gamma > 0$. An arbitrary but sufficiently weak local perturbation of H defines an $[[n, k, d - 2\ell, \delta(\ell)]]$ approximate error correcting code, where $\delta(\ell) \leq c_1 e^{-c_2 \ell / \log^2 \ell}$ for some constant $c_1, c_2 > 0$.*

Proof. The locality-preserving property of (47) can be cast into a more convenient form in the following. For any channel \mathcal{X}_M^M on a region M there exists a channel \mathcal{Y}_{M+r}^{M+r} such that

$$\sup_{\eta^{ABCR}} \mathfrak{B} \left(\mathcal{U}^\dagger \mathcal{X}^M \mathcal{U}(\eta^{ABCR}), \mathcal{Y}^{M+r}(\eta^{ABCR}) \right) \leq \epsilon, \quad (48)$$

where \mathcal{U} denotes the unitary conjugation channel by U , and η is any state, not necessarily a code state. This claim (48) with $\epsilon = \sqrt{2C e^{-cr / \log^2 r}}$ follows by setting $\mathcal{Y} = \mathcal{Y}^{M+r}$ from (47) and the triangle inequality of trace norm. Then, we can turn any correcting map for the unperturbed code space to a correcting map for the perturbed one, by sacrificing the locality and accuracy a little:

Lemma 17. *If a region A is (δ, ℓ) -correctable with respect to a code Π , then the region $Z = A^{-r}$ is $(\delta + 2\epsilon, \ell + 2r)$ -correctable with respect to the code $U \Pi U^\dagger$.*

The lemma is proved in Appendix E. □

We remark that the locality-preserving unitary U in the sense of Eq. (47) exists under a sole condition that there exists a gapped Hamiltonian path whose ground spaces are Π_0 and Π_1 [8]. Hence, by Lemma 17, the notion of locally correctable region is an invariant property of gapped phases of matter.

V. OUTLOOK

In this paper, we have introduced a coherent framework for analyzing local approximate quantum error correcting codes on lattices, which naturally includes weak local perturbations of commuting projector codes. Based only on very general properties of recovery maps on the code subspace, we have proved a tradeoff bound between n , k , d (Theorem 9), and given a constraints on the shape of logical operators (Theorem 13), resembling the results for commuting projector codes in Ref. [10, 11]. Furthermore, we have shown that if the logical operators of a code in 2D can be approximated by flexible string operators of constant width, then the degeneracy of the code subspace is constant, supporting one of the core assumptions of algebraic anyon theories.

We have assumed that there is one qubit (\mathbb{C}^2) per unit volume of the lattice. One can translate all our results to a situation where the local Hilbert space dimension is $q \geq 2$, simply by redefining the unit length as $1 \rightarrow (\log q)^{1/D}$. It is thus important for the last result on the degeneracy bound under the assumption of flexible strings that the local Hilbert space dimension is finite.

A number of important problems remain open, and we outline a few below.

Eastin and Knill [37] have proved that no exact error correcting code can have a universal set of transversal gates. This no-go result has been sharpened for local stabilizer codes on lattices by Bravyi and König [38], where it was shown that any transversal logical gate in a D -dimensional local stabilizer code is contained in the D th level of the Clifford hierarchy. In general, it is important to understand whether approximate error correcting codes would provide an avenue to local universal fault-tolerance in 2D or 3D that circumvents these no-go results.

As remarked in the Introduction, our notion of local recovery map is in part motivated by topological systems with anyons where pair-annihilation amounts to error recovery. We indeed have shown that the local recovery map exists for the ground space of a local commuting projector Hamiltonian and its perturbation. However, we have not explicitly discussed any larger class of phases of matter that do not necessarily have commuting Hamiltonian representatives, yet have a ground space that is a local approximate quantum error correcting code. One may wonder if, for any gapped Hamiltonian with stable gap against perturbations, the ground space can be regarded as a local approximate error correcting code. A more concrete problem would be to ask about the optimal relation between the parameters ℓ and δ under the local TQO condition for unfrustrated Hamiltonians [7].

Since our code is identified with a subspace of a physical Hilbert space, we have excluded subsystem codes [39] or more general subalgebra codes [40]. Also neglected is the thermal encoding, which has been studied in regards to self-correcting quantum memory [41–43]. Bény and Oreshkov [20, 40] have discussed information-disturbance tradeoffs for subalgebra codes, but it remains to be seen how these would apply with the local recovery maps. Extensions to these more general error correction schemes with approximate recovery maps might be relevant in connection to field theories with holographic duals [44, 45].

ACKNOWLEDGMENTS

STF was supported by the Australian Research Council via EQuS project number CE11001013 and by an Australian Research Council Future Fellowship FT130101744. JH was supported by the Pappalardo Fellowship in Physics while at MIT. MJK was supported by the Carlsberg fund and the Villum foundation. IK's research at Perimeter Institute was supported by the Government of Canada through Industry Canada and by the Province of Ontario through the Ministry of

-
- [1] Barbara M. Terhal, “Quantum error correction for quantum memories,” *Rev. Mod. Phys.* **87**, 307 (2015), [arXiv:1302.3428](#).
 - [2] Daniel Gottesman, “An introduction to quantum error correction and fault-tolerant quantum computation,” *Quantum Information Science and Its Contributions to Mathematics* **68**, 13–60 (2009), [0904.2557](#).
 - [3] Debbie W. Leung, M. A. Nielsen, Isaac L. Chuang, and Yoshihisa Yamamoto, “Approximate quantum error correction can lead to better codes,” *Phys. Rev. A* **56**, 2567–2573 (1997).
 - [4] Claude Crépeau, Daniel Gottesman, and Adam Smith, “Approximate quantum error-correcting codes and secret sharing schemes,” in *Advances in Cryptology – EUROCRYPT 2005: 24th Annual International Conference on the Theory and Applications of Cryptographic Techniques, Aarhus, Denmark, May 22–26, 2005. Proceedings*, edited by Ronald Cramer (Springer Berlin Heidelberg, Berlin, Heidelberg, 2005) pp. 285–301, [quant-ph/0503139](#).
 - [5] Gregory Moore and Nicholas Read, “Nonabelions in the fractional quantum hall effect,” *Nuclear Physics B* **360**, 362–396 (1991).
 - [6] Alexei Kitaev, “Anyons in an exactly solved model and beyond,” *Annals of Physics* **321**, 2–111 (2006).
 - [7] Spyridon Michalakis and Justyna P Zwolak, “Stability of frustration-free Hamiltonians,” *Communications in Mathematical Physics* **322**, 277–302 (2013).
 - [8] Matthew B Hastings and Xiao-Gang Wen, “Quasiadiabatic continuation of quantum states: The stability of topological ground-state degeneracy and emergent gauge invariance,” *Physical review b* **72**, 045141 (2005).
 - [9] Sergey Bravyi, Matthew B Hastings, and Spyridon Michalakis, “Topological quantum order: stability under local perturbations,” *Journal of mathematical physics* **51**, 093512 (2010).
 - [10] Sergey Bravyi, David Poulin, and Barbara Terhal, “Tradeoffs for reliable quantum information storage in 2D systems,” *Phys. Rev. Lett.* **104**, 050503 (2010), [arXiv:0909.5200](#).
 - [11] Jeongwan Haah and John Preskill, “Logical operator tradeoff for local quantum codes,” *Phys. Rev. A* **86**, 032308 (2012), [1011.3529](#).
 - [12] Sergey Bravyi and Barbara Terhal, “A no-go theorem for a two-dimensional self-correcting quantum memory based on stabilizer codes,” *New Journal of Physics* **11**, 043029 (2009).
 - [13] Nicolas Delfosse, “Tradeoffs for reliable quantum information storage in surface codes and color codes,” in *2013 IEEE International Symposium on Information Theory* (Institute of Electrical & Electronics Engineers (IEEE), 2013) [arXiv:1301.6588](#).
 - [14] Sergey Bravyi, “Subsystem codes with spatially local generators,” *Phys. Rev. A* **83**, 012320 (2011), [1008.1029](#).
 - [15] Dave Bacon, Steven T. Flammia, Aram W. Harrow, and Jonathan Shi, “Sparse Quantum Codes from Quantum Circuits,” in *Proceedings of the Forty-Seventh Annual ACM on Symposium on Theory of Computing*, STOC ’15 (ACM Press, New York, NY, USA, 2015) pp. 327–334, [arXiv:1411.3334](#).
 - [16] Carlo Cafaro and Peter van Loock, “Approximate quantum error correction for generalized amplitude-damping errors,” *Phys. Rev. A* **89** (2014), [10.1103/physreva.89.022316](#), [arXiv:1308.4582](#).
 - [17] Yingkai Ouyang, “Permutation-invariant quantum codes,” *Phys. Rev. A* **90** (2014), [10.1103/physreva.90.062317](#), [arXiv:1302.3247](#).
 - [18] Markus Grassl, Linghang Kong, Zhaohui Wei, Zhang-Qi Yin, and Bei Zeng, “Quantum Error-Correcting Codes for Qudit Amplitude Damping,” [arXiv:1509.06829](#) (2015), [arXiv:1509.06829](#).
 - [19] Benjamin Schumacher and Michael D. Westmoreland, “Approximate quantum error correction,” *Quant. Info. Process.* **1**, 5–12 (2002), [quant-ph/0112106](#).
 - [20] Cédric Bény and Ognian Oreshkov, “General conditions for approximate quantum error correction and near-optimal recovery channels,” *Phys. Rev. Lett.* **104**, 120501 (2010).
 - [21] Hui Khoo Ng and Prabha Mandayam, “Simple approach to approximate quantum error correction based on the transpose channel,” *Phys. Rev. A* **81**, 062342 (2010).
 - [22] Armin Uhlmann, “The “transition probability” in the state space of a $*$ -algebra,” *Reports on Mathematical Physics* **9**, 273–279 (1976).

- [23] C. A. Fuchs and J. van de Graaf, “Cryptographic distinguishability measures for quantum-mechanical states,” *IEEE Trans. Inf. Theory* **45**, 1216 (1999).
- [24] Emanuel Knill and Raymond Laflamme, “Theory of quantum error-correcting codes,” *Phys. Rev. A* **55**, 900–911 (1997).
- [25] Dennis Kretschmann, Dirk Schlingemann, and Reinhard F Werner, “The information-disturbance tradeoff and the continuity of Stinespring’s representation,” *IEEE transactions on information theory* **54**, 1708–1717 (2008).
- [26] Hui Khoo Ng and Prabha Mandayam, “Simple approach to approximate quantum error correction based on the transpose channel,” *Physical Review A* **81**, 062342 (2010).
- [27] Patrick Hayden and Andreas Winter, “Weak decoupling duality and quantum identification,” *IEEE Transactions on Information Theory* **58**, 4914–4929 (2012).
- [28] Benjamin Schumacher and M. A. Nielsen, “Quantum data processing and error correction,” *Phys. Rev. A* **54**, 2629–2635 (1996).
- [29] A. Yu. Kitaev, A. H. Shen, and M. N. Vyalyi, *Classical and Quantum Computation* (American Mathematical Society, 2002).
- [30] Martin Müller-Lennert, Frédéric Dupuis, Oleg Szehr, Serge Fehr, and Marco Tomamichel, “On quantum Rényi entropies: a new generalization and some properties,” *J. Math. Phys.* **54**, 122203 (2013), 1306.3142.
- [31] Salman Beigi, “Sandwiched Rényi divergence satisfies data processing inequality,” *J. Math. Phys.* **54**, 122202 (2013), 1306.5920.
- [32] Beni Yoshida and Isaac L. Chuang, “Framework for classifying logical operators in stabilizer codes,” *Phys. Rev. A* **81**, 052302 (2010), arXiv:1002.0085.
- [33] John Preskill, “Quantum error correction, Lecture notes for Physics 219, Caltech,” (1999).
- [34] Cumrun Vafa, “Toward classification of conformal theories,” *Physics Letters B* **206**, 421–426 (1988).
- [35] Matthew B Hastings, “Lieb-Schultz-Mattis in higher dimensions,” *Physical review b* **69**, 104431 (2004).
- [36] Sven Bachmann, Spyridon Michalakis, Bruno Nachtergaele, and Robert Sims, “Automorphic equivalence within gapped phases of quantum lattice systems,” *Communications in Mathematical Physics* **309**, 835–871 (2012).
- [37] Bryan Eastin and Emanuel Knill, “Restrictions on transversal encoded quantum gate sets,” *Physical review letters* **102**, 110502 (2009).
- [38] Sergey Bravyi and Robert König, “Classification of topologically protected gates for local stabilizer codes,” *Physical review letters* **110**, 170503 (2013).
- [39] David Poulin, “Stabilizer formalism for operator quantum error correction,” *Phys. Rev. Lett.* **95**, 230504 (2005), quant-ph/0508131.
- [40] Cédric Bény, “Conditions for the approximate correction of algebras,” in *Theory of Quantum Computation, Communication, and Cryptography* (Springer, 2009) pp. 66–75.
- [41] Eric Dennis, Alexei Kitaev, Andrew Landahl, and John Preskill, “Topological quantum memory,” *Journal of Mathematical Physics* **43**, 4452–4505 (2002).
- [42] Robert Alicki, Michał Horodecki, Paweł Horodecki, and Ryszard Horodecki, “On thermal stability of topological qubit in kitaev’s 4d model,” *Open Systems & Information Dynamics* **17**, 1–20 (2010).
- [43] Benjamin J Brown, Daniel Loss, Jiannis K Pachos, Chris N Self, and James R Wootton, “Quantum memories at finite temperature,” arXiv preprint arXiv:1411.6643 (2014).
- [44] Fernando Pastawski, Beni Yoshida, Daniel Harlow, and John Preskill, “Holographic quantum error-correcting codes: Toy models for the bulk/boundary correspondence,” arXiv preprint arXiv:1503.06237 (2015).
- [45] Daniel Harlow, “The Ryu-Takayanagi formula from quantum error correction,” 1607.03901v2.
- [46] Koenraad M R Audenaert, “A sharp continuity estimate for the von Neumann entropy,” *Journal of Physics A: Mathematical and Theoretical* **40**, 8127–8136 (2007).
- [47] R Alicki and M Fannes, “Continuity of quantum conditional information,” *J. Phys. A: Math. Gen.* **37**, L55–L57 (2004).

Appendix A: Proof of Theorem 3: Information-disturbance tradeoff

Theorem 18 (More general version of Theorem 3). *Let Π be a subspace on QC , and R a purifying space for Π . Given a Stinespring purification V^{QE} of a channel $\mathcal{N} : \rho \mapsto \text{Tr}_E(V^{QE}\rho V^{QE\dagger})$ for $\rho = \Pi\rho = \rho\Pi$, let $\mathcal{N}^c : \rho \mapsto \text{Tr}_Q(V^{QE}\rho V^{QE\dagger})$ be the complementary channel. Similarly, fix a purification W^{QE} of a channel \mathcal{M} on Q , and define the complementary channel \mathcal{M}^c . Then, we have*

$$\inf_{\mathcal{R}} \sup_{\rho^{QCR}} \mathfrak{B}(\mathcal{M}(\rho^{QCR}), \mathcal{R} \circ \mathcal{N}(\rho^{QCR})) = \inf_S \sup_{\rho^{QCR}} \mathfrak{B}(\mathcal{S} \circ \mathcal{M}^c(\rho^{QCR}), \mathcal{N}^c(\rho^{QCR})). \quad (\text{A1})$$

This theorem with the subsystem C being empty ($C = \mathbb{C}$) was stated in Ref. [20]. The proof in Ref. [20] omits the last step to replace an operator of norm at most 1 with a unitary operator [25]. Our statement is only different from that of Ref. [20] in that it includes a subsystem C explicitly. This enables us to accommodate locality. The domain and the codomain of the channels \mathcal{N}, \mathcal{M} being the same is mere convenience of presentation; more general cases reduce to the present one. The theorem implies in particular that when a code is defined on a physical system ABC ($Q = AB$), if $\mathcal{N} = \text{Tr}_A$, $\mathcal{N}^c = \text{Tr}_B$, $\mathcal{M} = \text{id}$, and \mathcal{M}^c outputs a fixed state ω , then

$$\inf_{\mathcal{R}_B^{AB}} \sup_{\rho^{ABCR}} \mathfrak{B}(\rho^{ABCR}, \mathcal{R}_B^{AB}(\rho^{BCR})) = \inf_{\omega^A} \sup_{\rho^{ABCR}} \mathfrak{B}(\omega^A \otimes \rho^{CR}, \rho^{ACR}), \quad (\text{A2})$$

which is Theorem 3.

Proof. Let $X^{QE'}$ denote a purification of the channel \mathcal{R} . The new environment E' is arbitrary here, unlike the subsystem E that is fixed by our choice of the complementary channel. Any channel (CPTP map) with d_i -dimensional input and d_o -dimensional output can be represented with a $(d_i d_o)$ -dimensional environment, which implies that the domain of all channels \mathcal{R} and \mathcal{S} is compact. However, for the present proof, it is useful to note that the optimizations over the channels \mathcal{R}, \mathcal{S} are with arbitrary environments.

Since the fidelity $\mathfrak{F} = 1 - \mathfrak{B}^2$ is concave in the arguments, the maximization over ρ^{QCR} can be restricted to pure states ρ^{QCR} . Assuming ρ^{QCR} is pure, let us express the fidelity on the left-hand side of Eq. (A1) using Uhlmann's theorem.

$$LHS = \sup_{X^{QE'}} \inf_{\rho^{QCR}} \sup_{Y^{EE'}} \langle \rho^{QCR} 0^{EE'} | W^{QE\dagger} Y^{EE'} X^{QE'} V^{QE} | \rho^{QCR} 0^{EE'} \rangle \quad (\text{A3})$$

$$= \sup_{X^{QE'}} \inf_{\rho^Q} \sup_{Y^{EE'}} \text{Tr} \left[\rho^Q 0^{EE'} W^{QE\dagger} Y^{EE'} X^{QE'} V^{QE} \right] \quad (\text{A4})$$

In the second line, the domain of ρ^Q is convex. (Indeed, $t\rho_1^Q + (1-t)\rho_2^Q$ is reduced from $t\rho_1^{QC} + (1-t)\rho_2^{QC}$, where $\rho_{1,2}^{QC}$ are some code states from the definition of $\rho_{1,2}^Q$, which can be purified using R .) Even if we relax the domain of $Y^{QE'}$ to those of the operator $Z^{QE'}$ of norm at most 1, the inner-most supremum always occurs when $Z^{QE'}$ is a unitary. This follows because $\sup_{U: \|U\| \leq 1} \Re \text{Tr} AU = \|A\|_1$ for any matrix A where the supremum is achieved by a unitary from the singular value decomposition of A . After such a relaxation, we can apply the minimax theorem to the bilinear objective function to obtain

$$LHS = \sup_{X^{QE'}} \sup_{Z^{EE'}} \inf_{\rho^Q} \text{Tr} \left[\rho^Q 0^{EE'} W^{QE\dagger} Z^{EE'} X^{QE'} V^{QE} \right] \quad (\text{A5})$$

$$= \sup_{X^{QE'}} \sup_{Z^{EE'}} \inf_{\rho^Q} \text{Tr} \left[\rho^Q 0^E W^{QE\dagger} \langle 0^{E'} | Z^{EE'} X^{QE'} | 0^{E'} \rangle V^{QE} \right]. \quad (\text{A6})$$

The second line is a rearrangement to make the linear operator $\langle 0^{E'} | Z^{EE'} X^{QE'} | 0^{E'} \rangle : QE \rightarrow QE$ more visible. Since $Z^{EE'}$ has norm at most 1, we can introduce another subsystem $E'' = \mathbb{C}^2$ and find a unitary $S^{EE'E''}$ such that $Z^{EE'} = \langle 0^{E''} | S^{EE'E''} | 0^{E''} \rangle$ by Lemma 20 below. Then we can replace $\langle 0^{E'} | Z^{EE'} X^{QE'} | 0^{E'} \rangle$ with $\langle 0^{E'E''} | S^{QE'E''} (X^{QE'} I^{E''}) | 0^{E'E''} \rangle$. Treating $E'E''$ as a larger environment E' , which was arbitrary, we see that the maximin value of (A6) with arbitrary E' can be achieved by unitary $Z^{EE'}$ and $X^{QE'}$.

The RHS of the theorem can be manipulated symmetrically with the role of $X^{QE'}$ and $Y^{EE'}$ interchanged, eventually leading to (A6) exactly. This completes the proof of the theorem. \square

Using the Fuchs-van de Graaf relation

$$\mathfrak{F}^2 + \mathfrak{T}^2 \leq 1 \leq \mathfrak{F} + \mathfrak{T},$$

we recover Theorem 3 of Ref. [25]:

Corollary 19. *For any (noise) channel \mathcal{N} , we have*

$$\frac{1}{4} \inf_{\mathcal{R}} \|\mathcal{R} \circ \mathcal{N} - \text{id}\|_{\diamond}^2 \leq \inf_{\mathcal{P}} \|\mathcal{N}^c - \mathcal{P}\|_{\diamond} \leq 2 \inf_{\mathcal{R}} \|\mathcal{R} \circ \mathcal{N} - \text{id}\|_{\diamond}^{1/2} \quad (\text{A7})$$

where \mathcal{P} denotes any completely depolarizing channel.

Lemma 20. *For any matrix M^A such that $\|M\| \leq 1$, there exists a unitary V^{AB} such that $\langle 0^B | V^{AB} | 0^B \rangle = M^A$ where $B = \mathbb{C}^2$.*

Proof. Let $M = XDY$ be a singular value decomposition of M with X, Y being unitary. If we had the lemma proved for the diagonal D , then multiplying $X^A \otimes I^B$ and $Y^A \otimes I^B$ will prove the lemma in the general case. Hence, it suffices to prove the lemma for any non-negative diagonal M .

By assumption, the non-negative diagonal elements of M are ≤ 1 . The operator V^{AB} has four blocks $\langle i^B | V^{AB} | j^B \rangle$ where $i, j = 0, 1$. Define $\langle 0^B | V^{AB} | 0^B \rangle = M$, and $\langle 1^B | V^{AB} | 0^B \rangle = \sqrt{1 - M^2}$. Then the columns of $V^{AB} | 0^B \rangle$ are orthonormal. There exists a basis of the vector space on AB that extends $V^{AB} | 0^B \rangle$, which can be used to fill the block $V^{AB} | 1^B \rangle$. \square

Appendix B: Proof of Theorem 5: Decoupling-Correctability

Lemma 21. *Two numbers*

$$\mu = \sup_{\rho^{ABCR}} \mathfrak{B}(\rho^{ACR}, \rho^A \otimes \rho^{CR}), \quad (\text{B1})$$

$$\delta = \inf_{\omega^A} \sup_{\rho^{ABCR}} \mathfrak{B}(\rho^{ACR}, \omega^A \otimes \rho^{CR}). \quad (\text{B2})$$

satisfy

$$\frac{1}{9} \delta^2 \leq \mu \leq 2\delta. \quad (\text{B3})$$

Proof. That $\mu \leq 2\delta$ follows because $\mathfrak{B}(\rho^A, \omega^A) \leq \delta$ by the monotonicity of the Bures distance. For the other inequality, we will show

$$\|\rho^A - \sigma^A\|_1 \leq 2(2 + \sqrt{2})\mu \quad (\text{B4})$$

for any code states ρ^{ABCR} and σ^{ABCR} . This implies that $\mathfrak{B}(\rho^A, \sigma^A) \leq \sqrt{(2 + \sqrt{2})\mu}$, and therefore

$$\delta \leq \sup_{\rho^{ABCR}} \mathfrak{B}(\rho^{ACR}, \sigma^A \otimes \rho^{CR}) \quad (\text{B5})$$

$$\leq \mathfrak{B}(\rho^{ACR}, \rho^A \otimes \rho^{CR}) + \sqrt{(2 + \sqrt{2})\mu} \quad (\text{B6})$$

$$\leq \left(1 + \sqrt{2 + \sqrt{2}}\right) \sqrt{\mu}. \quad (\text{B7})$$

Consider $\sup_{\rho, \sigma} \|\rho^A - \sigma^A\|_1$. Since for every ρ^A there is a code state ρ^{ABC} that reduces to ρ^A , we may regard the domain of optimization to be the set of mixed code states ρ^{ABC} . Since the trace distance is convex in its both arguments, the supremum occurs at extreme points, which are pure states. So, we only need to consider the case $|\rho^{ABC}\rangle = \sqrt{a}|\sigma^{ABC}\rangle + \sqrt{1-a}|\eta^{ABC}\rangle$ where $\langle \sigma^{ABC} | \eta^{ABC} \rangle = 0$ and $a \in [0, 1]$.

Define

$$|\psi^{ABCR}\rangle := \frac{1}{\sqrt{2}} (|\sigma^{ABC}0^R\rangle + |\eta^{ABC}1^R\rangle). \quad (\text{B8})$$

Then,

$$\Delta := \psi^{AR} - \psi^A \otimes \psi^R = \frac{1}{2} \begin{pmatrix} (\sigma^A - \eta^A)/2 & \underbrace{\text{Tr}_{BC} |\sigma^{ABC}\rangle \langle \eta^{ABC}|}_Y \\ \underbrace{\text{Tr}_{BC} |\eta^{ABC}\rangle \langle \sigma^{ABC}|}_{Y^\dagger} & -(\sigma^A - \eta^A)/2 \end{pmatrix}. \quad (\text{B9})$$

Since $\|X\|_1 = \sup_U \Re \text{Tr}(XU)$, we see

$$2\sqrt{2}\mu \geq \sup_U \Re \text{Tr} \left[\Delta \begin{pmatrix} U & 0 \\ 0 & -U \end{pmatrix} \right] \quad (\text{B10})$$

$$= \frac{1}{2} \sup_U \Re \text{Tr}[(\sigma^A - \eta^A)U] = \frac{1}{2} \|\sigma^A - \eta^A\|_1, \quad (\text{B11})$$

and also

$$2\sqrt{2}\mu \geq \sup_U \Re \text{Tr} \left[\Delta \begin{pmatrix} 0 & U \\ U & 0 \end{pmatrix} \right] = \frac{1}{2} \|Y + Y^\dagger\|_1. \quad (\text{B12})$$

Therefore,

$$\|\rho^A - \sigma^A\|_1 = \left\| (a-1)\sigma^A + (1-a)\eta^A + \sqrt{a(1-a)}(Y + Y^\dagger) \right\|_1 \quad (\text{B13})$$

$$\leq (1-a) \|\sigma^A - \eta^A\|_1 + \sqrt{a(1-a)} \|Y + Y^\dagger\|_1 \quad (\text{B14})$$

$$\leq 4\sqrt{2}(1-a)\mu + 4\sqrt{2}\sqrt{a(1-a)}\mu \quad (\text{B15})$$

$$\leq 2(2 + \sqrt{2})\mu. \quad (\text{B16})$$

□

Appendix C: Proof of Theorem 7: Cleaning Lemma and Correctability

Let R be a purifying system for the code space Π .

Lemma 22 (A local error correction map is a local cleaning map.). *Let A be a region of the physical system ABC . Suppose there exists a channel \mathcal{R}_B^{AB} such that*

$$\sup_{\rho^{ABC}} \|\mathcal{R}_B^{AB}(\rho^{BC}) - \rho^{ABC}\|_1 \leq \epsilon. \quad (\text{C1})$$

Then, for any operator U^{ABC} that is not necessarily logical, the pull-back $V^{BC} = (\mathcal{R}_B^{AB})^(U^{ABC})$ satisfies*

$$\|V^{BC}\| \leq \|U^{ABC}\| \quad (\text{C2})$$

$$\|\Pi(U^{ABC} - V^{BC})\Pi\| \leq 2\epsilon \|U^{ABC}\|. \quad (\text{C3})$$

In addition, if U^{ABC} is a logical unitary, then $V^{BC} = \mathcal{R}^(U^{ABC})$ satisfies*

$$\|V^{BC}\| \leq 1 \quad (\text{C4})$$

$$\|(U^{ABC} - V^{BC})\Pi\| \leq 4\sqrt{\epsilon} \quad (\text{C5})$$

$$\|(U^{ABC\dagger} - V^{BC\dagger})\Pi\| \leq 4\sqrt{\epsilon} \quad (\text{C6})$$

where Π is the projector onto the code space.

Proof. Since $V^{BC} = (W_B^{AB})^\dagger U^{ABC} W_B^{AB}$ for some Stinespring isometry W_B^{AB} , it follows that $\|V^{BC}\| \leq \|U^{ABC}\|$. Whenever U^{ABC} is hermitian, V^{BC} is also hermitian. Since $V^{BC\dagger} = (\mathcal{R}_B^{AB})^*(U^{ABC\dagger})$, (C6) will follow from (C5).

First, we estimate the norm restricted to the code space assuming U^{ABC} is hermitian.

$$\|\Pi(U^{ABC} - V^{BC})\Pi\| = \sup_{\rho^{ABC}} |\text{Tr}(\rho^{ABC}(U^{ABC} - V^{BC}))| \quad (\text{C7})$$

$$= \sup_{\rho^{ABC}} |\text{Tr}(\rho^{ABC}U^{ABC} - \mathcal{R}_B^{AB}(\rho^{BC})U^{ABC})| \quad (\text{C8})$$

$$= \sup_{\rho^{ABC}} |\text{Tr}((\rho^{ABC} - \mathcal{R}_B^{AB}(\rho^{BC}))U^{ABC})| \quad (\text{C9})$$

$$\leq \sup_{\rho^{ABC}} \|\rho^{ABC} - \mathcal{R}_B^{AB}(\rho^{BC})\|_1 \|U^{ABC}\| \quad (\text{C10})$$

$$\leq \epsilon \|U^{ABC}\|. \quad (\text{C11})$$

For general U^{ABC} , decompose U^{ABC} into hermitian and anti-hermitian parts, and use the triangle inequality for the operator norm. This proves Eq. (C3).

To prove Eq. (C5), let $|\psi\rangle = \Pi|\psi\rangle$ be any normalized code state. If $\alpha = \|\Pi V^{BC}|\psi\rangle\|_2$ and $\beta = \|(I - \Pi)V^{BC}|\psi\rangle\|_2$, then $\alpha \geq 1 - 2\epsilon$ by Eq. (C3) and the unitarity of U^{ABC} . Since $\|V^{BC}\| \leq \|U^{ABC}\| = 1$, we see $\alpha^2 + \beta^2 \leq 1$. Hence, $\beta^2 \leq 1 - \alpha^2 \leq 1 - (1 - 2\epsilon)^2 \leq 4\epsilon$, and

$$\|(I - \Pi)V^{BC}\Pi\| \leq 2\sqrt{\epsilon}. \quad (\text{C12})$$

Therefore,

$$\|(U^{ABC} - V^{BC})\Pi\| \leq \|(I - \Pi)(U^{ABC} - V^{BC})\Pi\| + \|\Pi(U^{ABC} - V^{BC})\Pi\| \quad (\text{C13})$$

$$\leq 2\sqrt{\epsilon} + 2\epsilon \quad (\text{C14})$$

$$\leq 4\sqrt{\epsilon} \quad (\text{C15})$$

□

Lemma 23 (A cleanable region is correctable.). *Suppose for any unitary U^{AB} that preserves the code space there exists an operator V^B supported on B such that*

$$\|V^B\| \leq 1, \quad (\text{C16})$$

$$\|(U^{AB} - V^B)\Pi\| \leq \epsilon, \quad (\text{C17})$$

$$\|(U^{AB\dagger} - V^{B\dagger})\Pi\| \leq \epsilon. \quad (\text{C18})$$

Then, for any code state ρ^{ABR} with purifying space R

$$\|\rho^{AR} - \omega^A \otimes \rho^R\|_1 \leq 5\epsilon. \quad (\text{C19})$$

Proof. Using the triangle inequality and the fact that $\|PQ\|_1, \|QP\|_1 \leq \|P\|_1 \cdot \|Q\|$, we have

$$\|U^{AB} \rho^{ABR} U^{AB\dagger} - V^B \rho^{ABR} V^{B\dagger}\|_1 \leq \|U^{AB} \rho^{ABR} (U^{AB\dagger} - V^{B\dagger})\|_1 \quad (\text{C20})$$

$$+ \|(U^{AB} - V^B) \rho^{ABR} V^{B\dagger}\|_1 \quad (\text{C21})$$

$$\leq 2\epsilon, \quad (\text{C22})$$

$$\|V^{B\dagger} V^B \Pi - \Pi\| = \|V^{B\dagger} \Pi V^B \Pi - U^{AB\dagger} \Pi U^{AB} \Pi + V^{B\dagger} (I - \Pi) V^B \Pi\| \quad (\text{C23})$$

$$\leq \|(V^{B\dagger} - U^{AB\dagger}) \Pi V^B \Pi + U^{AB\dagger} \Pi (V^B - U^{AB}) \Pi\| \quad (\text{C24})$$

$$+ \|V^{B\dagger} (I - \Pi) (V^B - U^{AB}) \Pi\| \quad (\text{C25})$$

$$\leq 3\epsilon. \quad (\text{C26})$$

Since a partial trace cannot increase the trace distance, we see

$$\|\rho^{AR} - \text{Tr}_B (U^{AB} \rho^{ABR} U^{AB\dagger})\|_1 \leq \|\rho^{AR} - \text{Tr}_B (V^B \rho^{ABR} V^{B\dagger})\|_1 + 2\epsilon \quad (\text{C27})$$

$$= \|\rho^{AR} - \text{Tr}_B (V^{B\dagger} V^B \Pi \rho^{ABR})\|_1 + 2\epsilon \quad (\text{C28})$$

$$\leq \|\rho^{AR} - \text{Tr}_B (\rho^{ABR})\|_1 + 5\epsilon \quad (\text{C29})$$

$$= 5\epsilon \quad (\text{C30})$$

where U^{AB} is an arbitrary unitary that preserves the code space. A completely depolarizing channel over the code space can be written as

$$\int dU^{AB} U^{AB} \rho^{ABR} U^{AB\dagger} = \rho^R \otimes \Pi / \text{Tr}(\Pi) \quad (\text{C31})$$

where the integral is over the group of all logical unitaries. Using Eq. (C30), we see that

$$\left\| \int dU^{AB} \left[\rho^{AR} - \text{Tr}_B (U^{AB} \rho^{ABR} U^{AB\dagger}) \right] \right\|_1 \leq \int dU^{AB} \left\| \rho^{AR} - \text{Tr}_B (U^{AB} \rho^{ABR} U^{AB\dagger}) \right\|_1 \quad (\text{C32})$$

$$\leq 5\epsilon. \quad (\text{C33})$$

Let $\omega^A = \text{Tr}_B \Pi / \text{Tr}(\Pi)$. Then,

$$\|\rho^{AR} - \omega^A \otimes \rho^R\|_1 = \left\| \int dU^{AB} \rho^{AR} - \omega^A \otimes \rho^R \right\|_1 \quad (\text{C34})$$

$$\leq \left\| \int dU^{AB} \text{Tr}_B (U^{AB} \rho^{ABR} U^{AB\dagger}) - \omega^A \otimes \rho^R \right\|_1 + 5\epsilon \quad (\text{C35})$$

$$= 5\epsilon. \quad (\text{C36})$$

□

Appendix D: Proof of Theorem 4: Disentangling Unitary

Lemma 24. *Let Π be a code space on ABC , and R be a purifying space. Let $\mathcal{V}_B^{B'B''} = V \cdot V^\dagger$ denote any isometry channel, where B' is some auxiliary system and B'' is a copy of AB . Then,*

$$\inf_{\omega^A} \sup_{\rho^{ABCR}} \mathfrak{B}(\rho^{ACR}, \omega^A \otimes \rho^{CR}) = \inf_{\omega^{AB'}, \mathcal{V}_B^{B'B''}} \sup_{\rho^{ABCR}} \mathfrak{B}(\mathcal{V}_B^{B'B''}(\rho^{ABCR}), \omega^{AB'} \otimes \rho^{B''CR}) \quad (\text{D1})$$

where $\rho^{B''CR}$ is the same as ρ^{ABCR} but supported on B'' instead of AB .

Proof of 24 as a corollary of the information-disturbance theorem. That $LHS \leq RHS$ follows at once because the Bures distance can only decrease by tracing out $B'B''$.

For the other direction, we need to find an isometry and a state $\omega^{AB'}$ that yields the desired minimax value. Since \mathfrak{B}^2 is convex in its arguments, the supremum over ρ^{ABCR} always occurs at a pure state, and thus we can only consider when ρ^{ABCR} is pure. The information-disturbance theorem (Thm. 3) provides a recovery map \mathcal{R}_B^{AB} to fidelity $1 - \delta^2$ where $\delta = LHS$. Let $V = V_B^{ABE}$ be a Stinespring isometry of \mathcal{R}_B^{AB} . The fidelity being at least $1 - \delta^2$ means that

$$|\langle \omega^{DE} \rho^{ABCR} | V_B^{ABE} | \rho^{DBCR} \rangle| \geq 1 - \delta^2 \quad (\text{D2})$$

for some pure state ω^{DE} where ρ^{DBCR} is the same as ρ^{ABCR} but supported on $DBCR$ instead of $ABCR$. Interchanging the subsystem label D and A , we have

$$|\langle \omega^{AE} \rho^{DBCR} | V_B^{DBE} | \rho^{ABCR} \rangle| \geq 1 - \delta^2. \quad (\text{D3})$$

Now we set $B' = E$ and $B'' = DB$ to finish the proof. \square

Conversely, the recovery map with respect to the noise Tr_A in the information-disturbance theorem can be constructed using the disentangling isometry $\mathcal{V}_B^{B'B''}$. Suppose the initial pure state is ρ^{ABCR} . The noise erases A , and we are confronted with ρ^{BCR} . Applying \mathcal{V} , we have, to fidelity $\geq 1 - \delta^2$, $\omega^{B'} \otimes \rho^{B''CR}$. Discarding B' and mapping B'' to AB isomorphically, the recovery succeeds with fidelity $\geq 1 - \delta^2$.

Direct Proof of 24. That $LHS \leq RHS$ follows at once because the Bures distance can only decrease by tracing out $B'B''$.

To show $\delta := LHS \geq RHS$, fix an ω^A that achieves the LHS in Eq. (D1).

$$\inf_{\rho^{ABCR}} \mathfrak{F}(\rho^{ACR}, \omega^A \otimes \rho^{CR}) = 1 - \delta^2 \quad (\text{D4})$$

Since the fidelity is concave in its arguments, the infimum occurs at a pure state ρ^{ABCR} . Introduce purification auxiliary systems D and EF where EF is a copy of AB . We apply Uhlmann's theorem with purifications $\omega^{AD} 0^B$ of ω^A and ρ^{EFCR} of ρ^{CR} . Define $S = S^\dagger$ to be the swap operator between AB and EF . Then we have,

$$1 - \delta^2 = \inf_{\text{pure } \rho^{ABCR}} \sup_{W^{BDEF}} \Re \langle \omega^{AD} 0^B \rho^{EFCR} | W^{BDEF} | \rho^{ABCR} 0^{DEF} \rangle \quad (\text{D5})$$

$$= \inf_{\text{pure } \rho^{ABCR}} \sup_{W^{BDEF}} \Re \langle \omega^{ED} 0^F \rho^{ABCR} | S^{ABEF} W^{BDEF} | \rho^{ABCR} 0^{DEF} \rangle \quad (\text{D6})$$

$$= \inf_{\rho^{ABC}} \sup_{W^{BDEF}} \Re \text{Tr} [S^{ABEF} W^{BDEF} \rho^{ABC} | 0^{DEF} \rangle \langle \omega^{ED} 0^F |] . \quad (\text{D7})$$

In the last expression, the domain of ρ^{ABC} is the set of all code states, which is convex. The domain of W^{BDEF} is the set of all unitaries, but can be relaxed to the set of all operators U on $BDEF$

of operator norm ≤ 1 , which is convex, since the supremum occurs at a unitary. The minimax theorem can then be applied since the objective function is linear in both arguments.

Reversing the manipulation, we have

$$1 - \delta^2 = \sup_{U^{BDEF} \text{ pure}} \inf_{\rho^{ABCR}} \Re \langle \omega^{AD} 0^B \rho^{EFCR} | U^{BDEF} | \rho^{ABCR} 0^{DEF} \rangle. \quad (\text{D8})$$

Here the optimal operator U is not necessarily a unitary. However, the supremum value can be reached by a unitary because for any operator U with $\|U\| \leq 1$ there is a unitary V such that $U^{BDEF} = \langle 0^{D'} | V^{BDD'EF} | 0^{D'} \rangle$ due to Lemma 20. The latter amounts to a larger choice of the purifying space D for ω^A . Hence, for this choice of V , the fidelity between two states on *RHS* of Eq. (D1) is at least $1 - \delta^2$. This is to say that $RHS \leq \delta = LHS$. Therefore, the claim is proved with $V_B^{B'B''} = V^{BDEF} | 0^{DEF} \rangle$ where $B' = BD$ and $B'' = EF$. \square

Appendix E: Proof of Lemma 17

Proof. Let \mathcal{R}_B^{AB} be a recovery map for Tr_A :

$$\sup_{\rho^{ABCR}} \mathfrak{B}(\mathcal{R}_B^{AB}(\rho^{BCR}), \rho^{ABCR}) = \delta. \quad (\text{E1})$$

Let $B_i = A \cap B^{+r}$ and $B_o = C \cap B^{+r}$, so that $B^{+r} = B_i B B_o$. Also, let $Z = A \setminus B^{+r} = A \setminus B_i$. We will show that there exists a recovery channel $\mathcal{T} = \mathcal{T}_{B^{+r}}^{ZB^{+r}}$ for the code $U\Pi U^\dagger$ such that

$$\sup_{\sigma^{ZB_i BCR} \in U\Pi U^\dagger \otimes R} \mathfrak{B}(\mathcal{T}(\sigma^{B_i BCR}), \sigma^{ZB_i BCR}) \leq 2\epsilon + \delta. \quad (\text{E2})$$

Although the domain of \mathcal{R}_B^{AB} is the set of all states supported on B , we may regard the channel \mathcal{R}_B^{AB} as the composition of the complete depolarization \mathcal{E}_A^A on A , followed by a channel \mathcal{D}_{AB}^{AB} . That is, $\mathcal{R}_B^{AB}(\sigma^B) = \mathcal{D}_{AB}^{AB} \circ \mathcal{E}_A^A(\sigma^B)$ for *any* state σ^B that partial-traces to σ^B . We will use this trick in what follows.

We wish to recover from the erasure of the “interior” $Z = A \setminus B^{+r}$ for the code $U\Pi U^\dagger$. We consider the following channel \mathcal{S} , which is going to be a recovery map,

$$\mathcal{S}(\sigma^{B_i BCR}) = \mathcal{U} \circ \mathcal{D} \circ \mathcal{U}^\dagger(I^Z/d_Z \otimes \sigma^{B_i BCR}). \quad (\text{E3})$$

We first calculate the fidelity of the recovery by \mathcal{S} . Let $\sigma^{ZB_i BCR} = \mathcal{U}(\rho^{ABCR})$ be an arbitrary code state in $U\Pi U^\dagger$, possibly entangled with reference system R . The recovery is successful up to error

$$\mathfrak{B}(\mathcal{S}(\sigma^{B_i BCR}), \sigma^{ZB_i BCR}) \quad (\text{E4})$$

$$= \mathfrak{B}(\mathcal{U}^\dagger \circ \mathcal{S} \circ \mathcal{E}_Z^Z \circ \mathcal{U}(\rho^{ZB_i BCR}), \rho^{ZB_i BCR}) \quad (\text{E5})$$

$$= \mathfrak{B}(\mathcal{D} \mathcal{U}^\dagger \mathcal{E}_Z^Z \mathcal{U}(\rho^{ZB_i BCR}), \rho^{ZB_i BCR}). \quad (\text{E6})$$

U is locality-preserving, and the channel $\mathcal{U}^\dagger \mathcal{E}_Z^Z \mathcal{U}$ can be replaced by a channel \mathcal{E}' on $Z^{+r} = A$ at the cost of additional error at most ϵ . Hence, we can bound the error as

$$\mathfrak{B}(\mathcal{S}(\sigma^{B_i BCR}), \sigma^{ZB_i BCR}) \leq \epsilon + \mathfrak{B}(\mathcal{D} \circ \mathcal{E}'(\rho^{ABCR}), \rho^{ABCR}) \quad (\text{E7})$$

$$= \epsilon + \delta \quad (\text{E8})$$

where the last step is because $\text{Tr}_A \mathcal{E}'(\rho^{ABCR}) = \rho^{BCR}$.

Now we investigate the support of \mathcal{S} . By assumption on U , the composition $\mathcal{U} \circ \mathcal{D} \circ \mathcal{U}^\dagger$ can be approximated by a channel \mathcal{T}' supported on $(AB)^{+r} = ABB_o$ to an error ϵ . Since the domain of \mathcal{S} is restricted to the states on $B_i BCR$, we may say that there exists $\mathcal{T}_{B_i B B_o}^{ZB_i B B_o}$ that recovers the erasure of Z for the code $U\Pi U^\dagger$ to error $2\epsilon + \delta$. \square

Appendix F: Continuity of mutual information

The Fannes-Audeanaert inequality [46] reads

$$|S(\rho) - S(\sigma)| \leq H(t) + t \log(d - 1) \quad (\text{F1})$$

for any d -dimensional states ρ and σ , where $t = \frac{1}{2} \|\rho - \sigma\|_1$, and $H(t) = -t \log t - (1 - t) \log(1 - t)$ is the binary entropy function.

Alicki and Fannes [47] have proved that when $t = \frac{1}{2} \|\rho^{AB} - \sigma^{AB}\|_1 < \frac{1}{2}$,

$$|S(\rho^{AB}) - S(\rho^B) - S(\sigma^{AB}) + S(\sigma^B)| \leq 2H(2t) + 8t \log d_A \quad (\text{F2})$$

where d_A is the dimension of the subsystem A .

Assume that subsystems A and B are nontrivial. It follows that the difference of the mutual information is bounded as

$$|I_\rho(A : B) - I_\sigma(A : B)| = |S(\rho^A) - S(\sigma^A) - S(\rho^{A|B}) - S(\sigma^{A|B})| \quad (\text{F3})$$

$$\leq H(t) + t \log(d_A - 1) + 2H(2t) + 8t \log d_A \quad (\text{F4})$$

$$\leq 9t \log d_A + 6t \log(1/t) \quad (\text{F5})$$

$$\leq 9t \log(d_A/t). \quad (\text{F6})$$

Since the mutual information is symmetric, d_A can be replaced by $\min(d_A, d_B) \geq 2$.